

Engines of Parsimony

computing on a budget

Hannah Earley

hannah.ie · DAMTP, Cambridge · Vaire Computing

physics of computing

computational performance constraints

dissipation in reversible computing (RC)

super-adiabatic RC?

parallelism & concurrency in RC

Thermodynamics of Computing

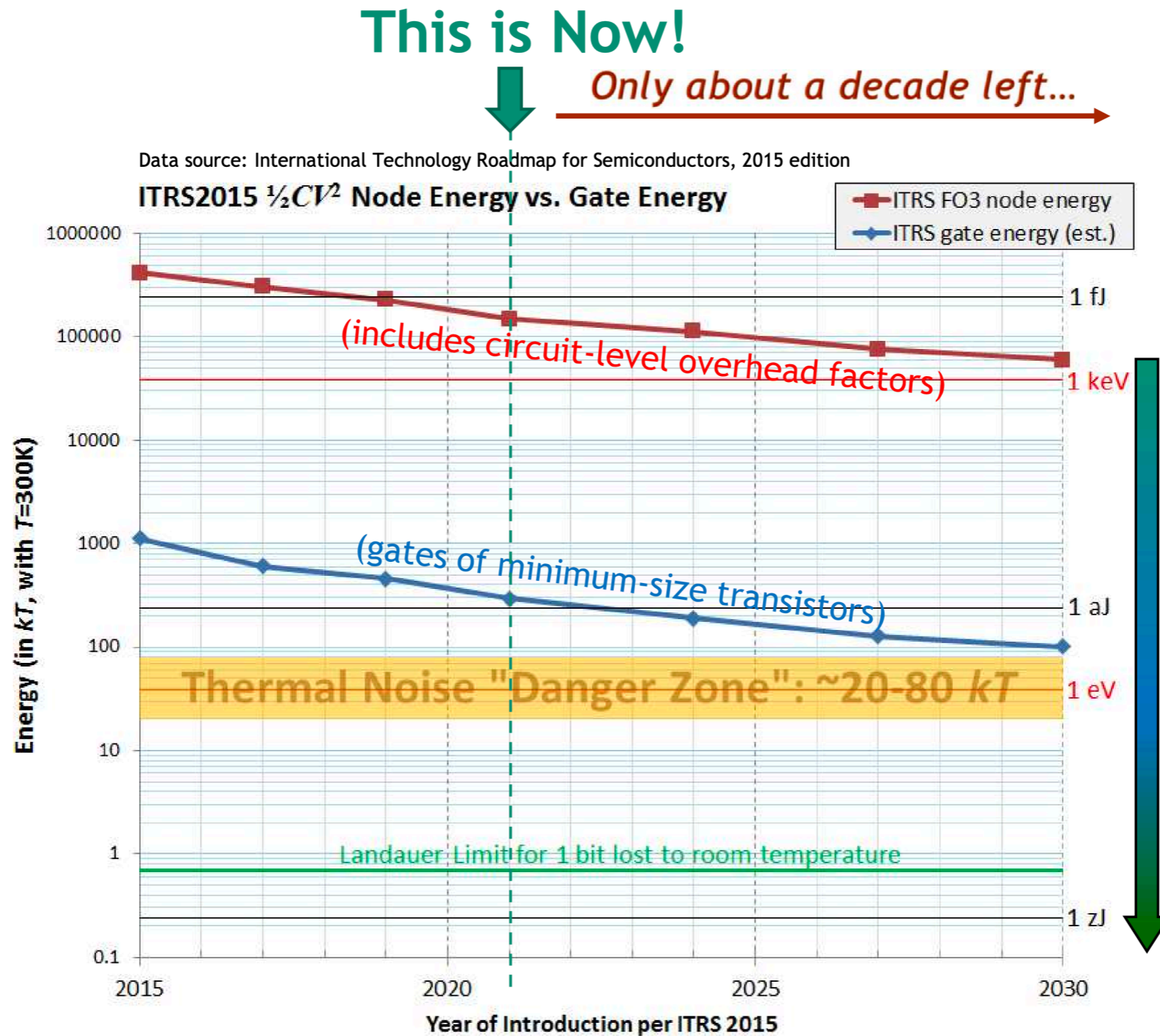
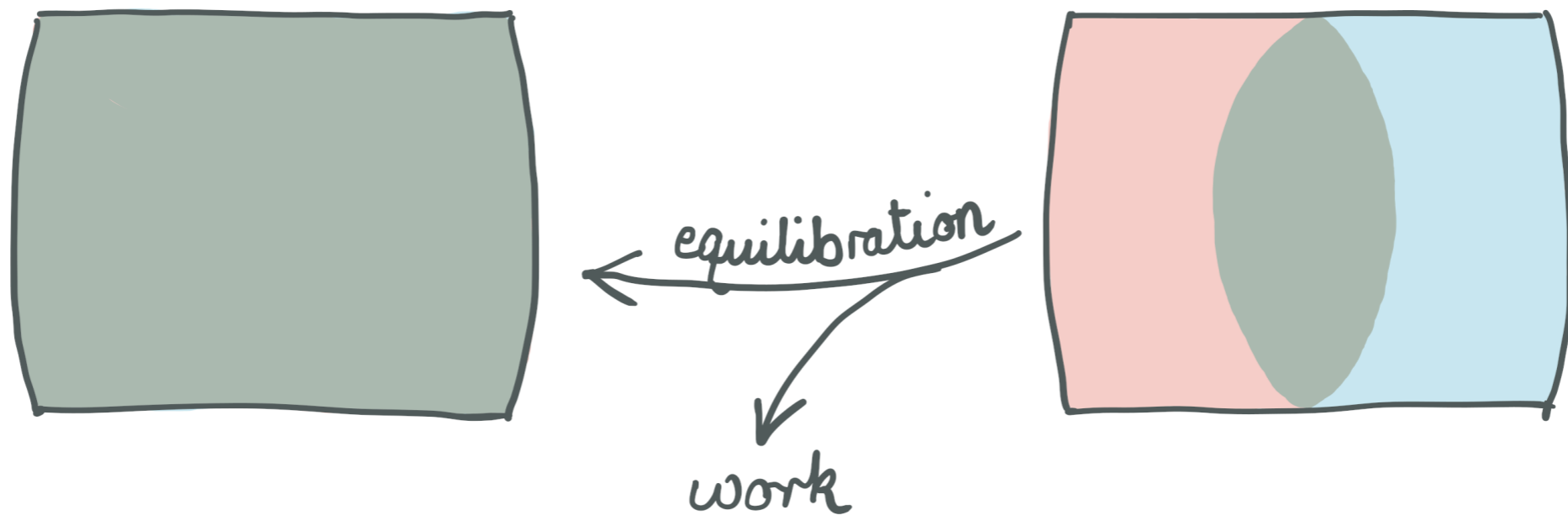
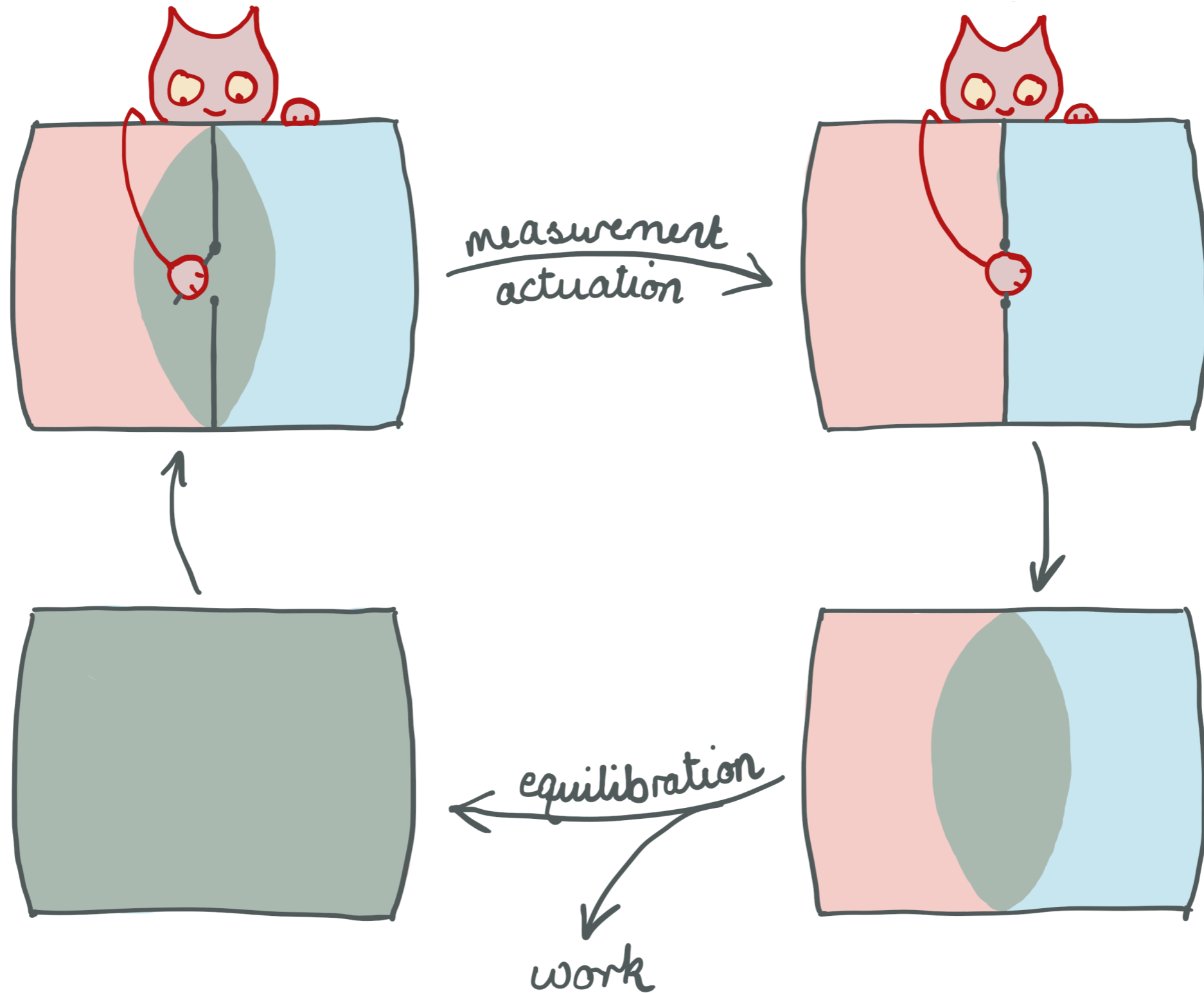


Figure – Mike Frank

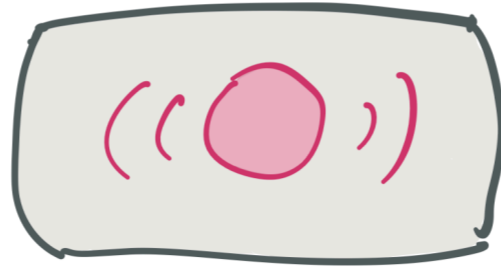
Maxwell's Daemon



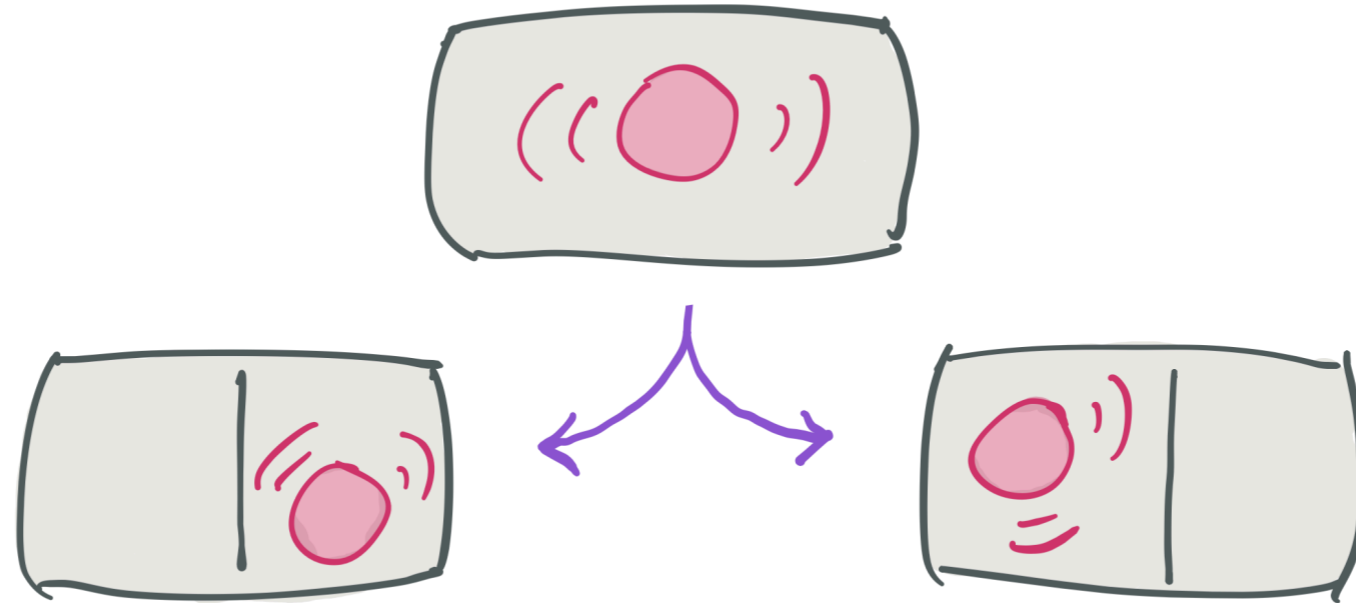
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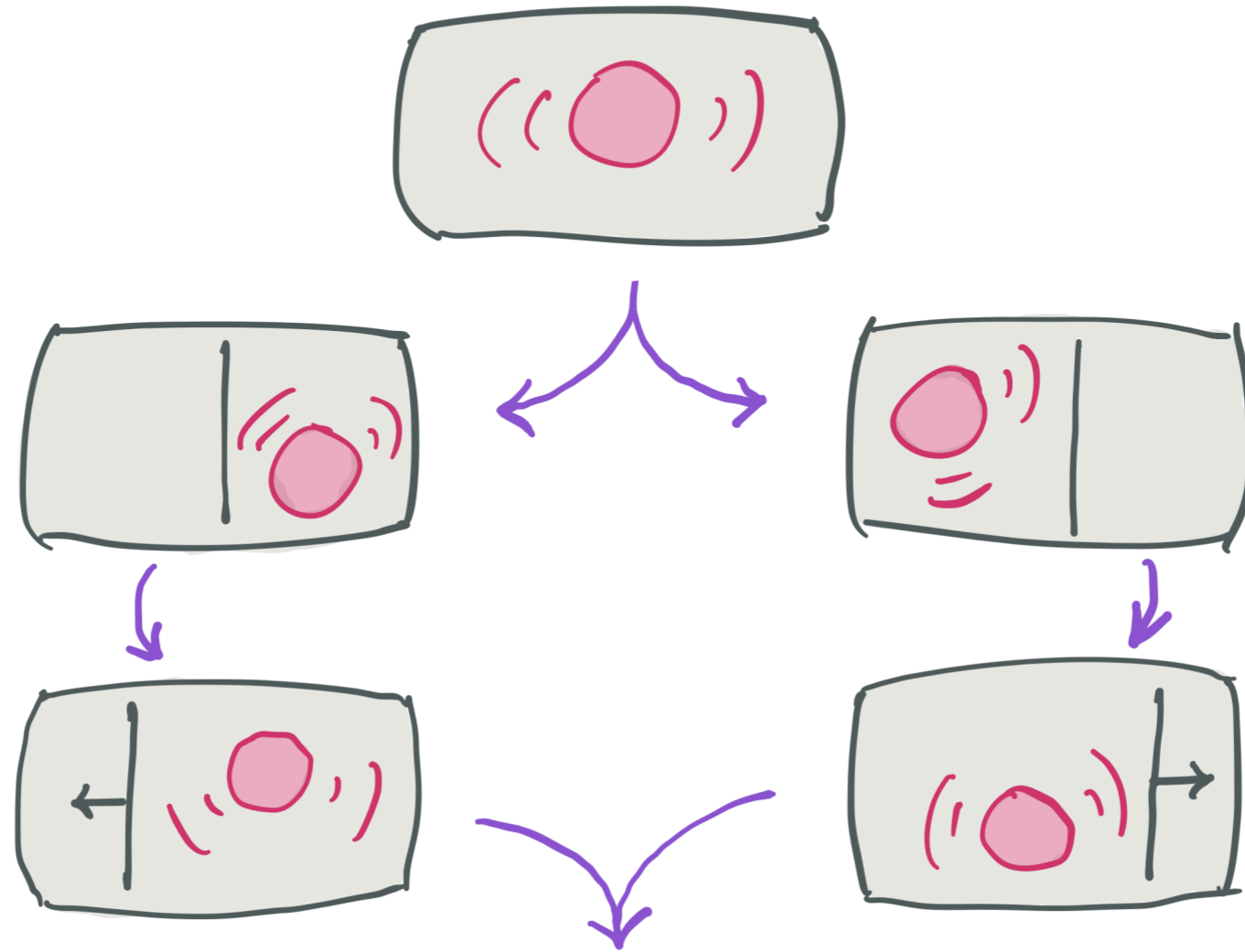
Szilard Engine



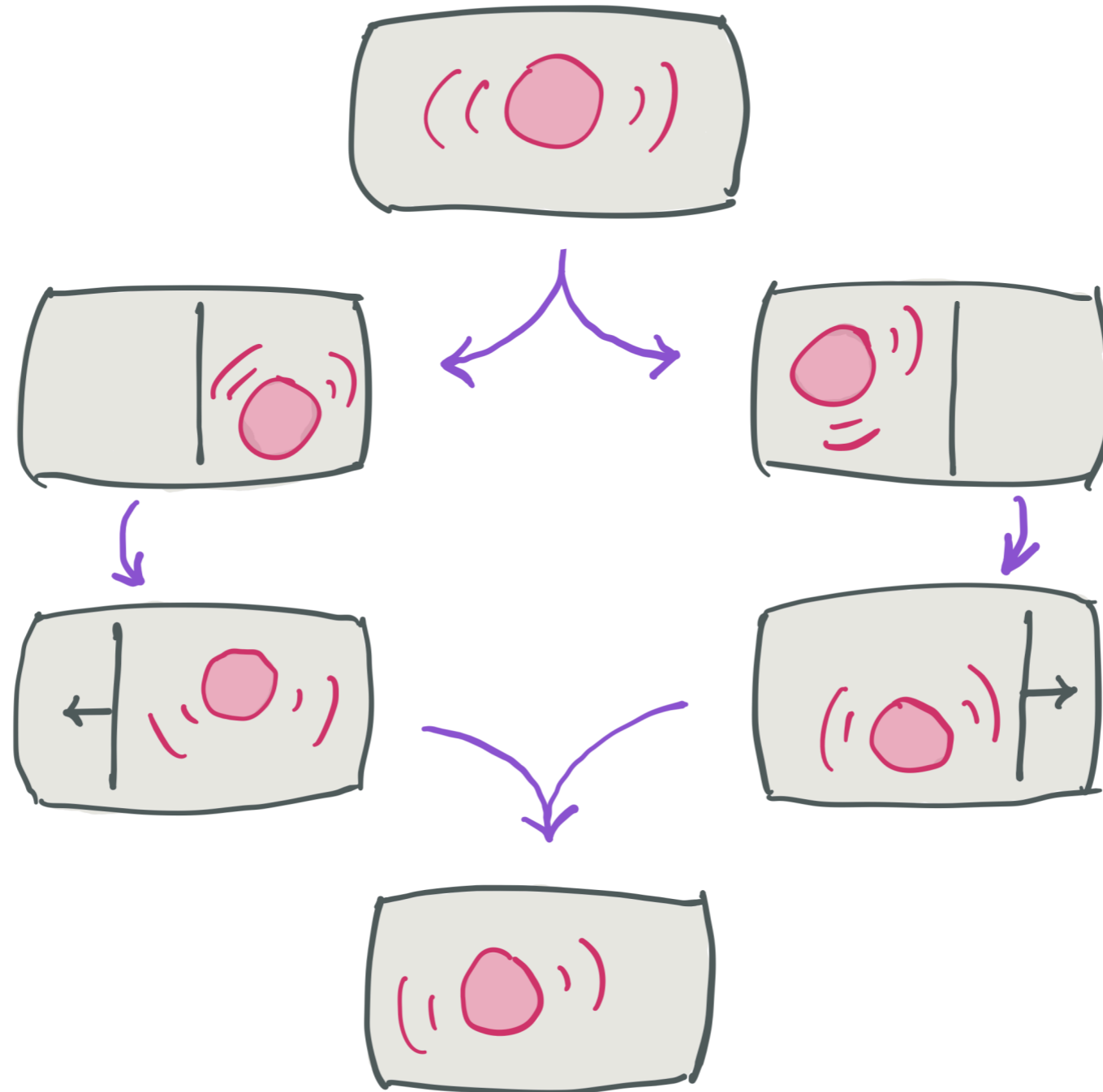
Szilard Engine



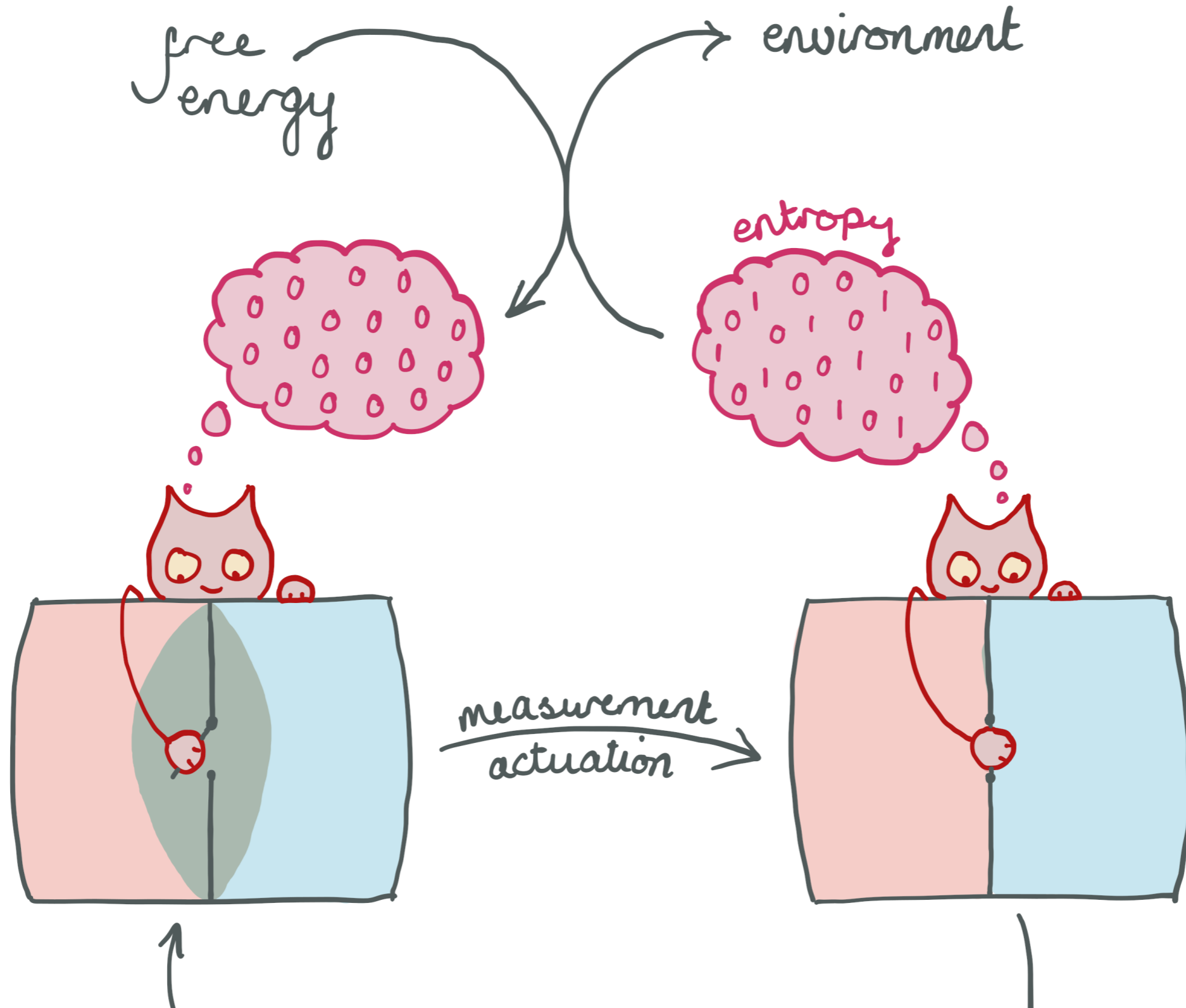
Szilard Engine



Szilard Engine



Maxwell's Daemon



Landauer Limit

$$\Delta S \geq k_B \Delta I$$

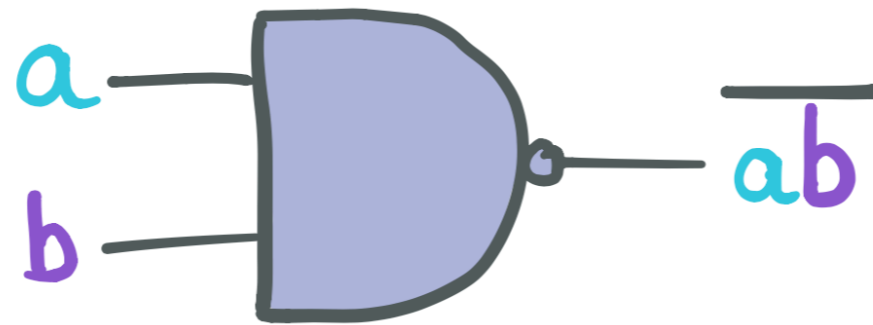
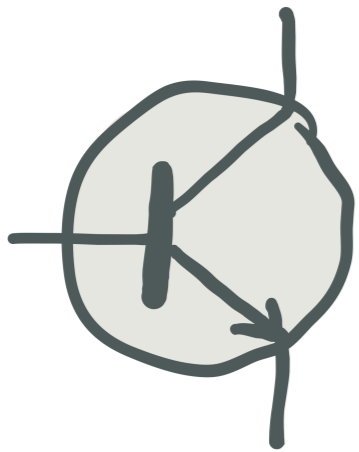
$$\Delta Q \geq k_B T \log 2$$

Landauer Limit

$$\Delta S \geq \Delta I$$

$$\Delta Q \geq T \log 2$$

Consequences for Computing

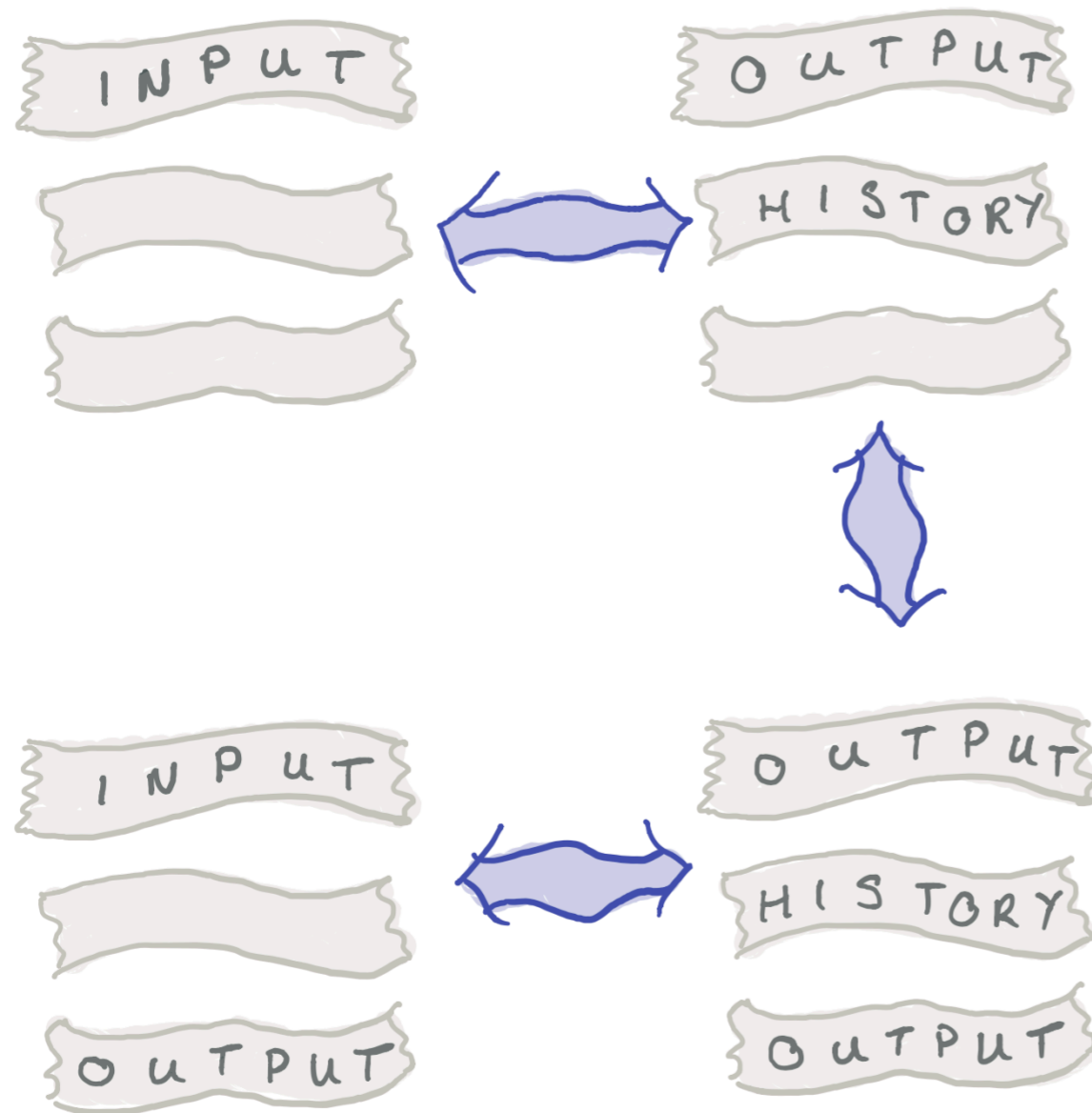


```
mov AL, 314  
jne loop
```

```
void swap(int *a, int *b) {  
    int tmp = *b;  
    *b = *a;  
    *a = tmp;  
}
```

```
int i = 5;  
while (i--) {  
    printf("%d\n", i);  
}
```

Reversible Computing


 $n \ x[2]$

procedure *fib*

if $n = 0$

then $x_0 += 1$

$x_1 += 1$

else $n -= 1$

call *fib*

$x_0 += x_1$

$x_0 \Leftrightarrow x_1$

fi $x_0 = x_1$

Bennett, Charles H. "Logical reversibility of computation." IBM journal of Research and Development 17.6 (1973): 525-532.

Yokoyama, Tetsuo, and Robert Glück. "A reversible programming language and its invertible self-interpreter." Proceedings of the 2007 ACM SIGPLAN symposium on Partial evaluation and semantics-based program manipulation. 2007.

physics of computing

computational performance constraints

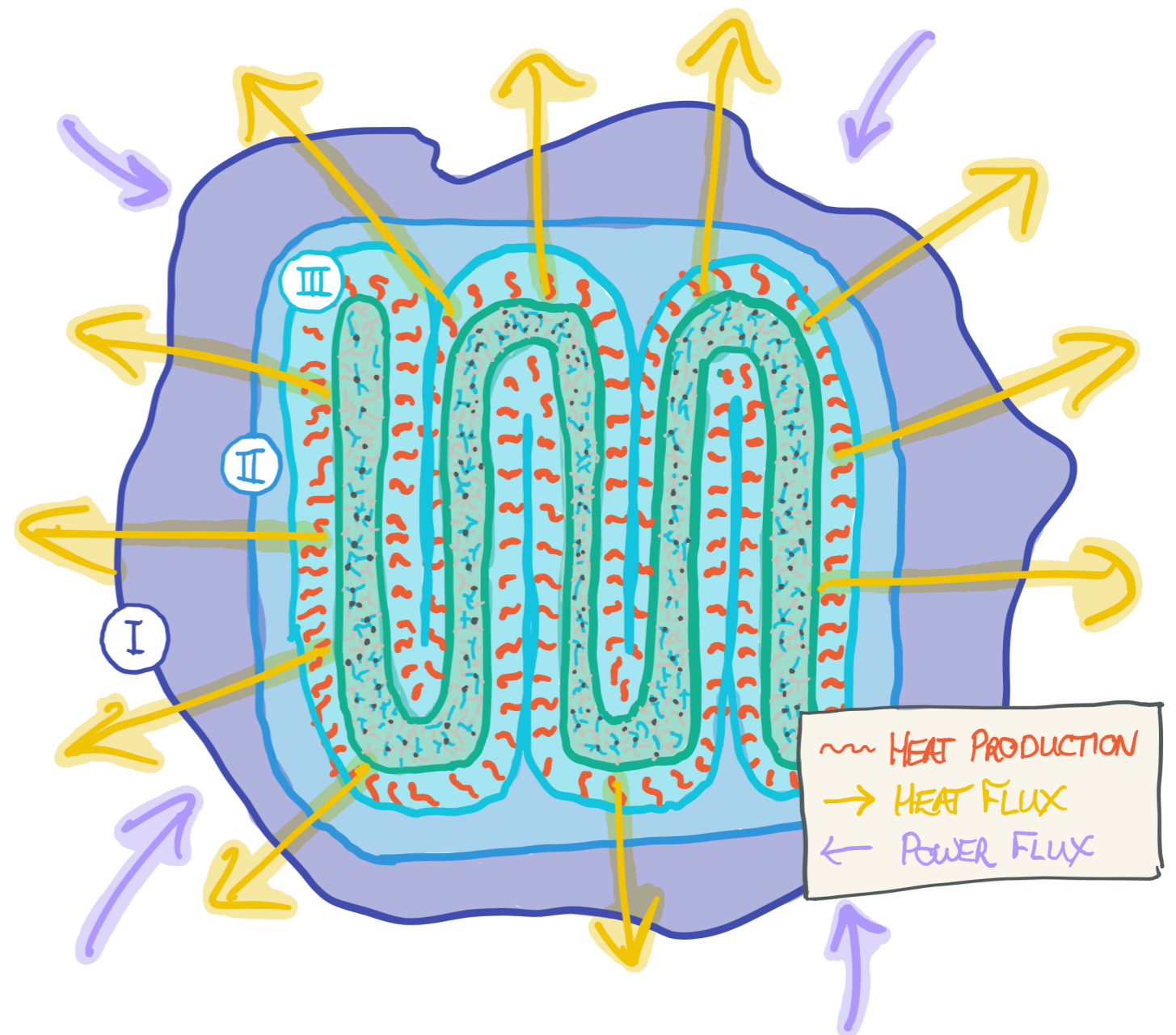
dissipation in reversible computing (RC)

super-adiabatic RC?

parallelism & concurrency in RC

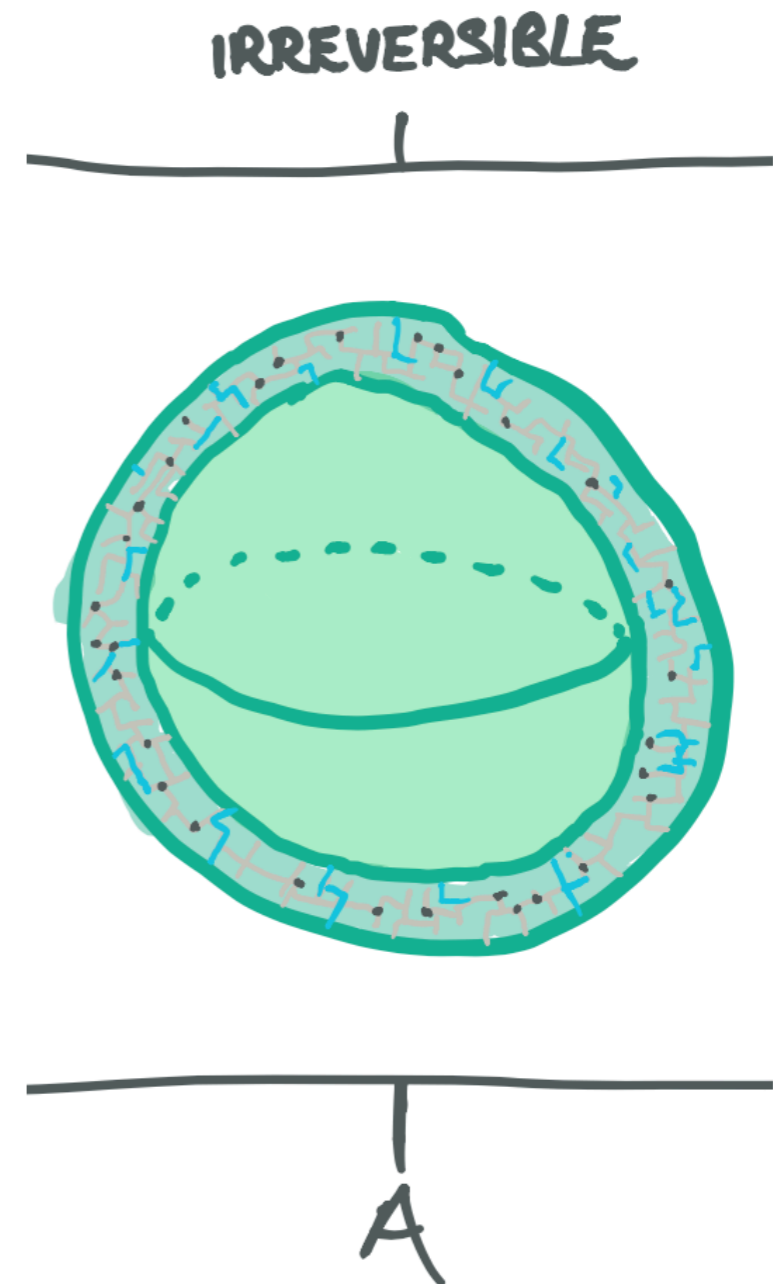
Setup

- mass-energy available for computation
- prop to volume
- heat/free energy flux
- prop to convex boundary



Conventional Computers

- Landauer Limit lower bound
- P/A constrained

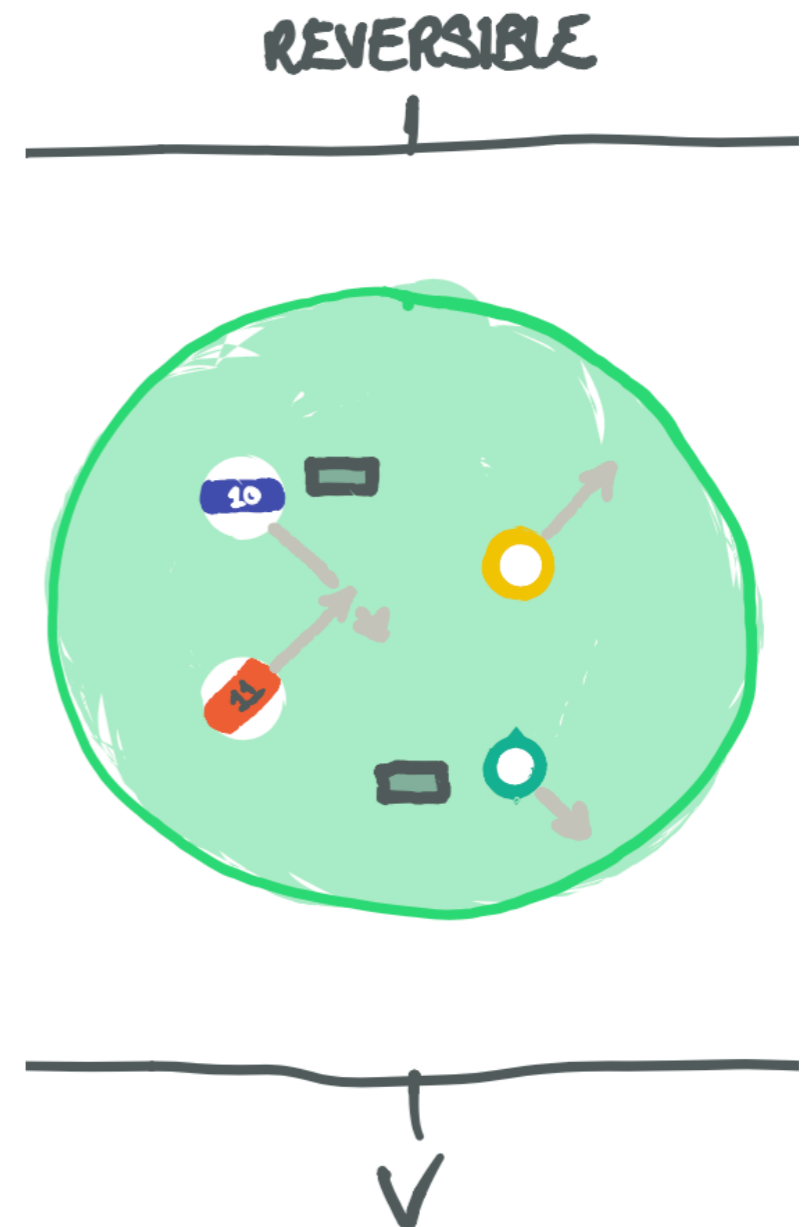


Ballistic Computing

- RC not subject to Landauer limit
- Computing for free(...?)
- Not P-limited...
- ...but M-limited

$$\nu \leq 2E/h$$

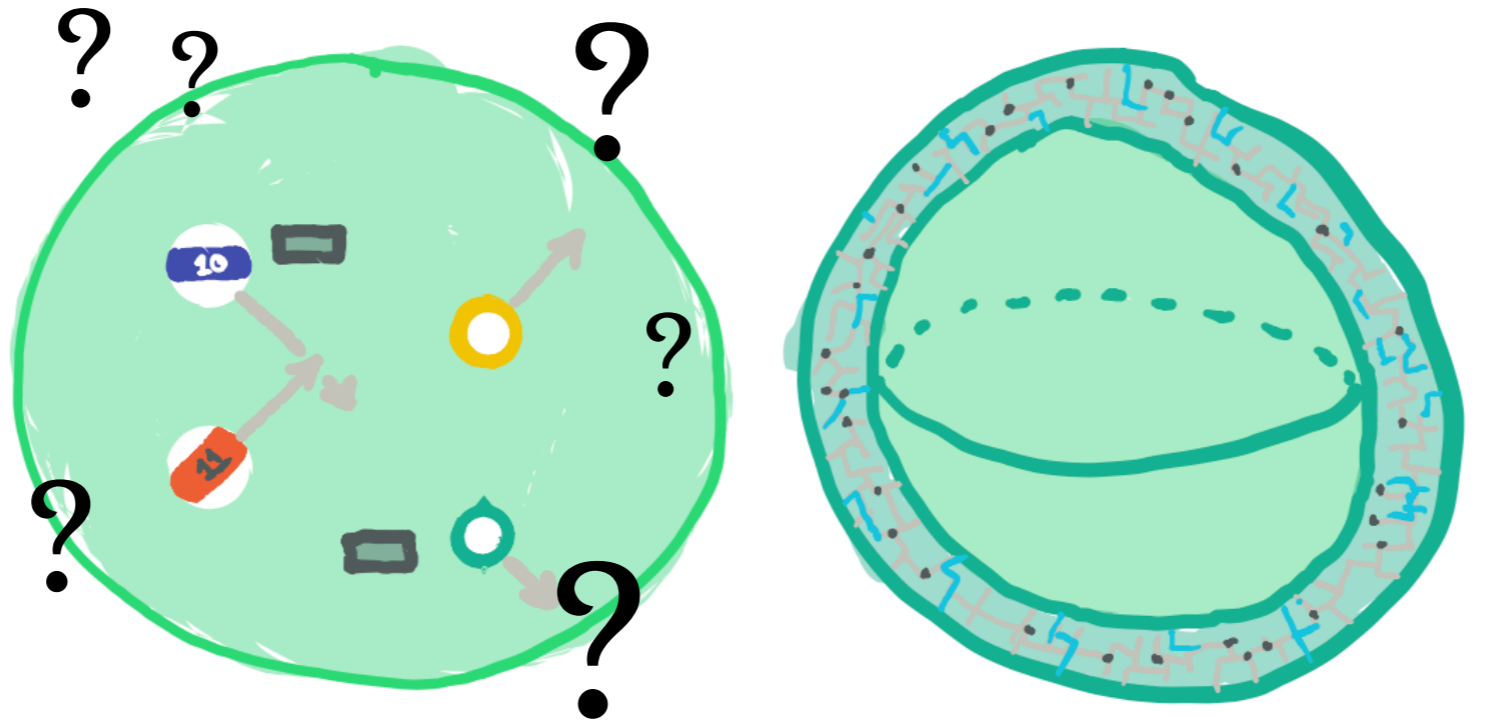
$$\leq E/\pi$$



architecture

REVERSIBLE

IRREVERSIBLE



V A

net rate of computation

physics of computing

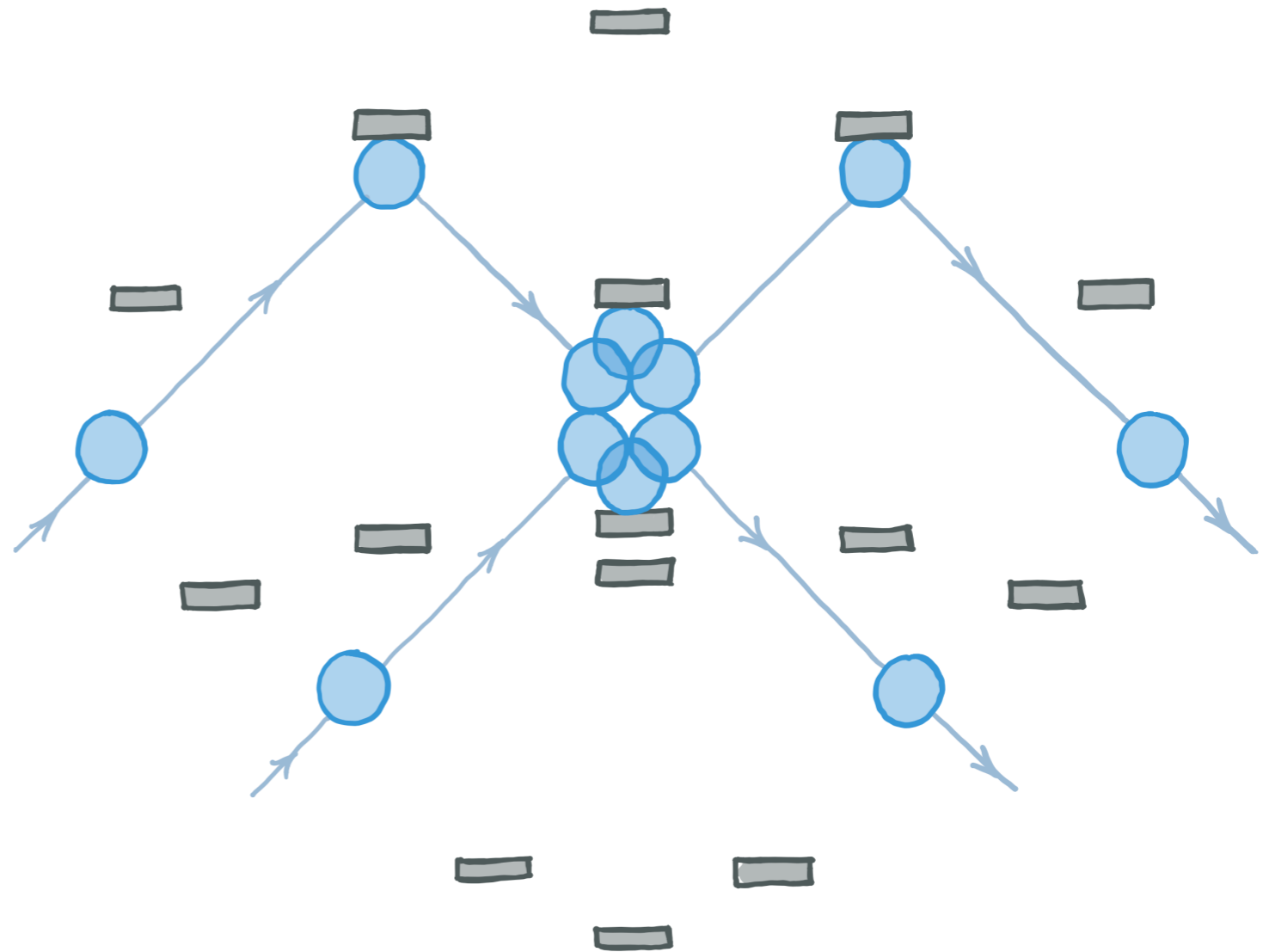
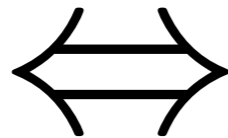
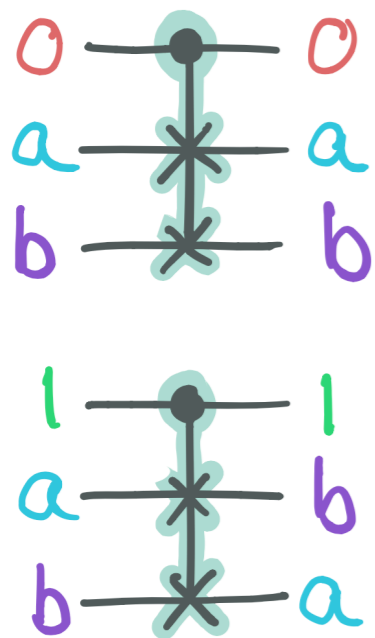
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Ballistic Computing?



Ballistic Computing?

“Even if classical balls could be shot with perfect accuracy into a perfect apparatus, fluctuating tidal forces from turbulence in the atmospheres of nearby stars would be enough to randomize their motion within a few hundred collisions. Needless to say, the trajectory would be spoiled much sooner if stronger nearby noise sources (e.g., thermal radiation and conduction) were not eliminated.”

-Charles Bennett

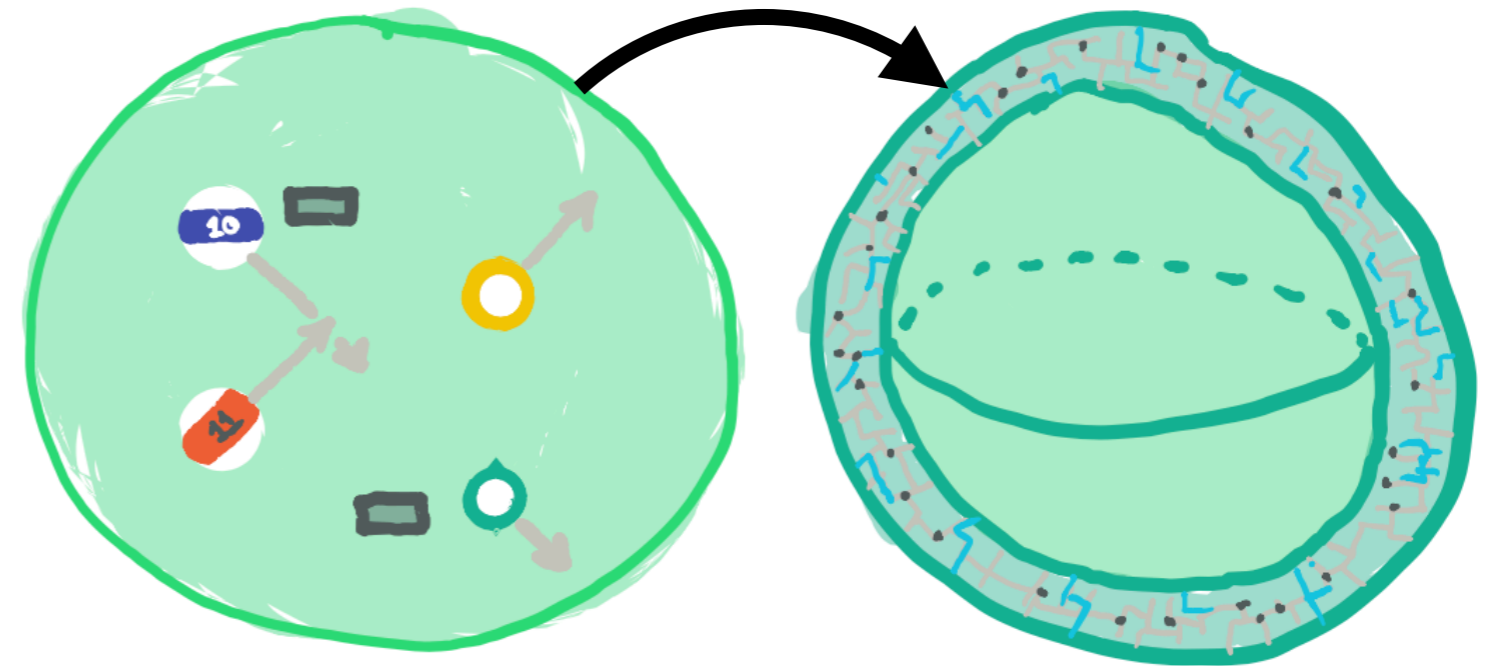
Dissipation in RC

architecture

REVERSIBLE

IRREVERSIBLE

?



- thermo coupling
- signal/state restoration
- superconductivity?
only DC...

✓

A

net rate of computation

Adiabatic RC

- asymptotically vanishing (but non-zero) dissipation

$$\Delta S \cdot \Delta \tau = \text{const.}$$

$$\Delta S \propto \nu$$

- realizable! adiabatic CMOS, ...

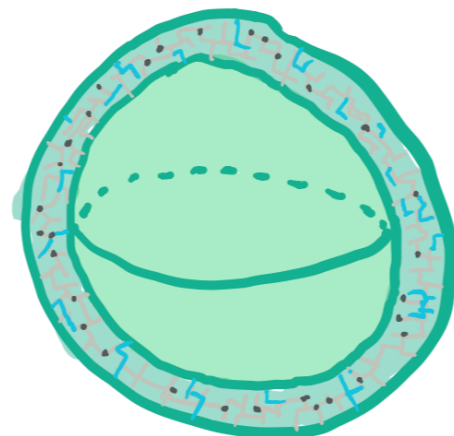
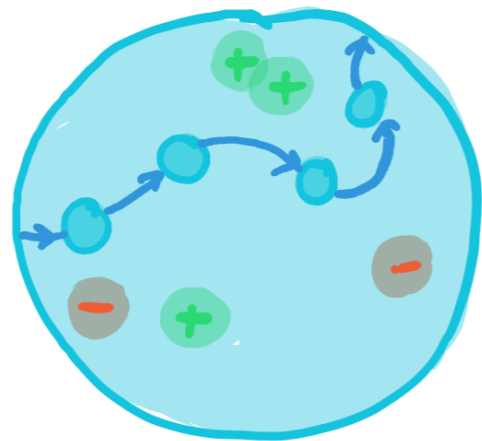
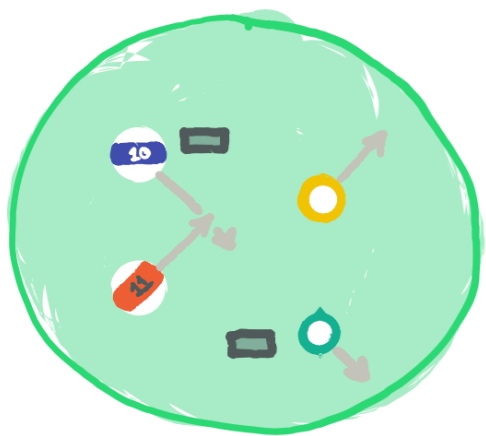
Adiabatic RC

architecture

REVERSIBLE

REVERSIBLE

IRREVERSIBLE



V

\sqrt{VA}

A

net rate of computation

Adiabatic RC

architecture

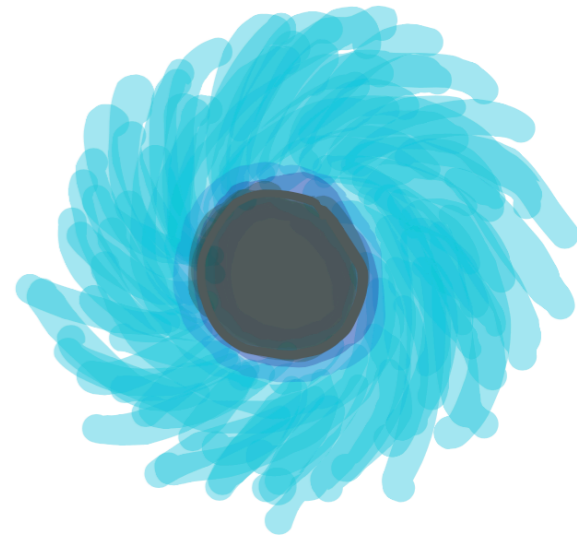
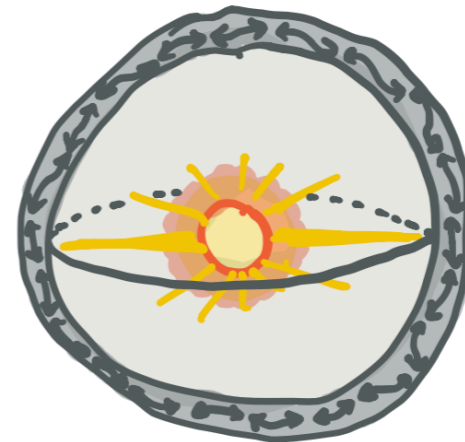
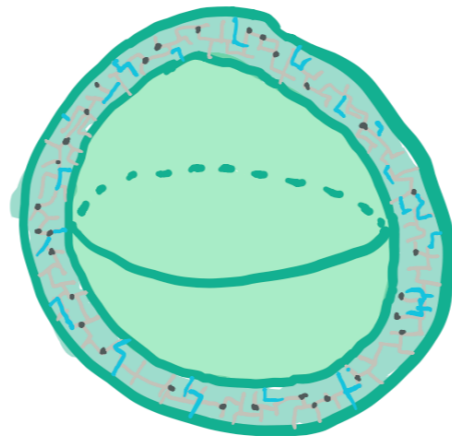
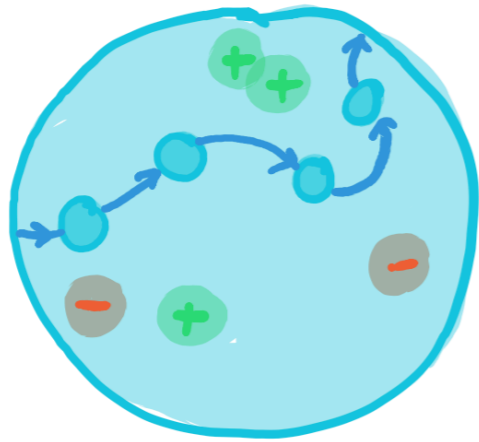
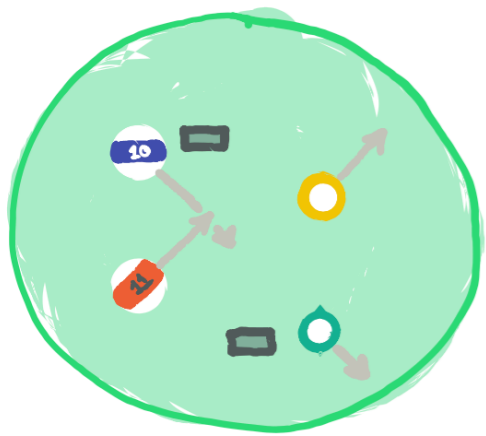
REVERSIBLE

REVERSIBLE

IRREVERSIBLE

REVERSIBLE

BOTH



V

\sqrt{VA}

A

\sqrt{AR}

R

net rate of computation

physics of computing

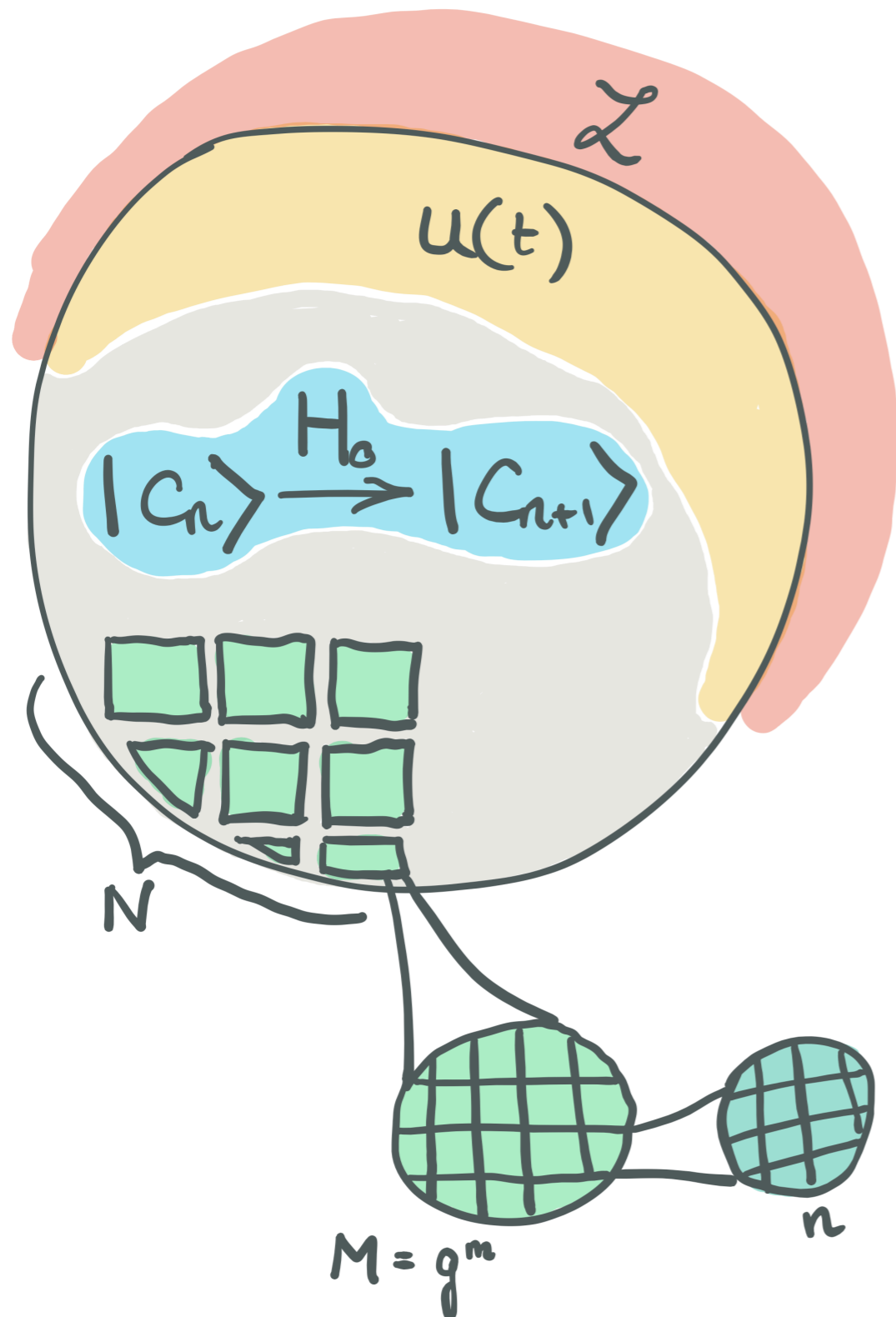
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Generic RC Model

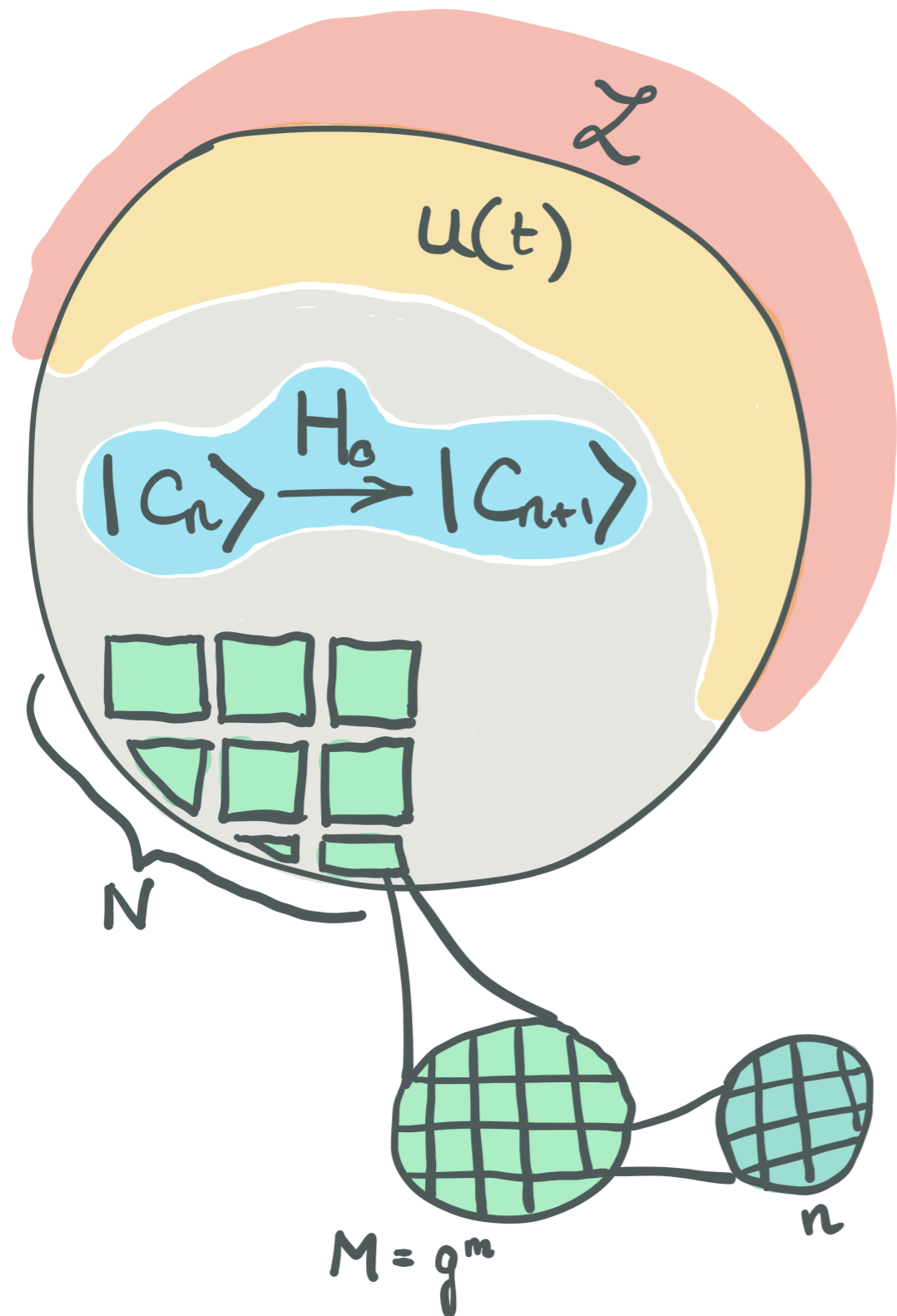


$$H = H_0 + U(t)$$

$$\dot{\rho} = i[\rho, H] + \mathcal{L}\rho$$

$$\mathcal{L}\rho = \sum_i \gamma_i (L_i \rho L_i^\dagger - \frac{1}{2} \{L_i^\dagger L_i, \rho\})$$

Generic RC Model



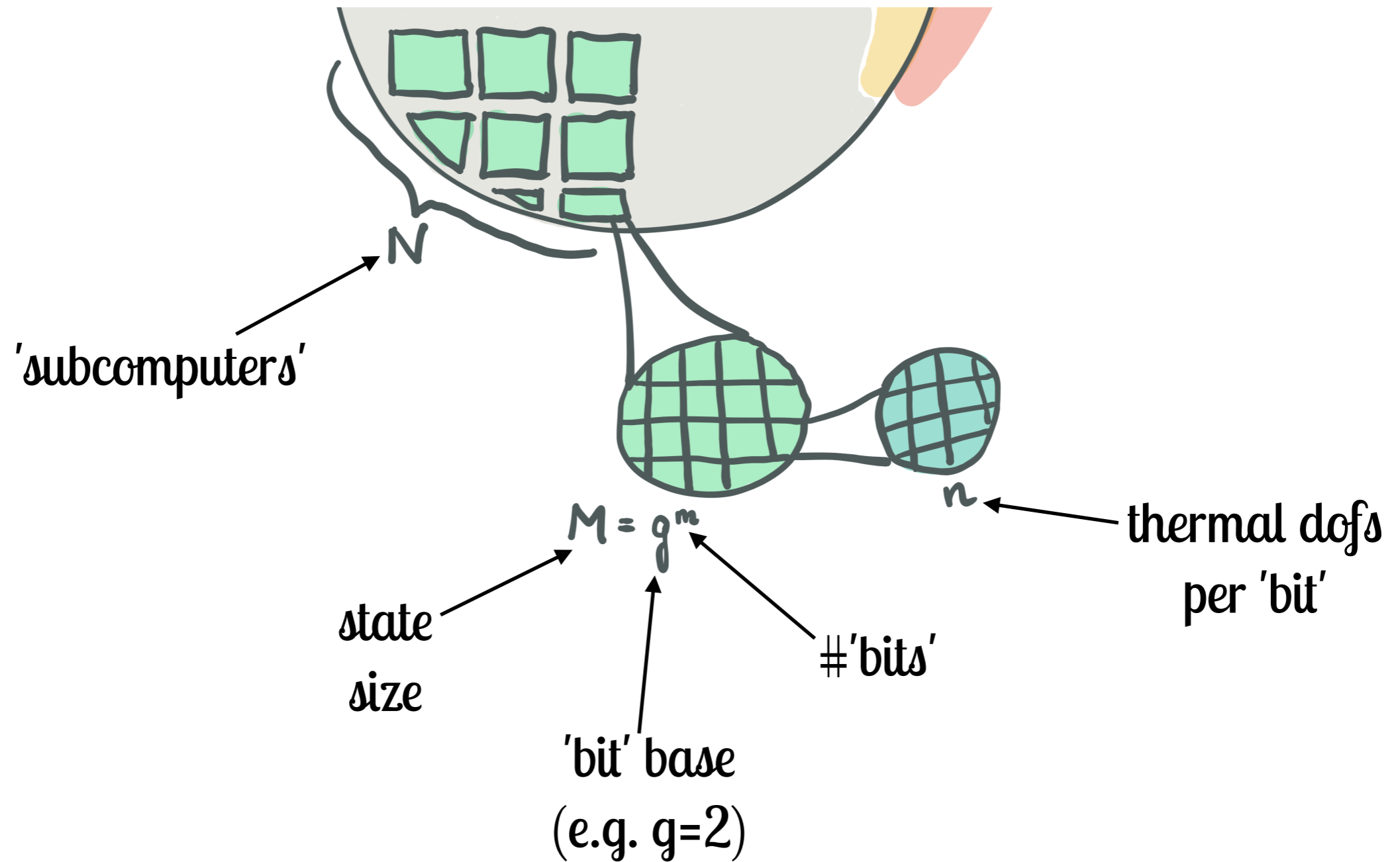
$$H = \overset{\text{ballistic}}{\downarrow} H_0 + \overset{\text{internal temp}}{\swarrow} U(t)$$

$$\dot{\rho} = i[\rho, H] + \overset{\text{external env}}{\uparrow} \mathcal{L}\rho$$

$$\mathcal{L}\rho = \sum_i \overset{\text{strength}}{\downarrow} \gamma_i (L_i \rho L_i^\dagger - \frac{1}{2} \{L_i^\dagger L_i, \rho\})$$

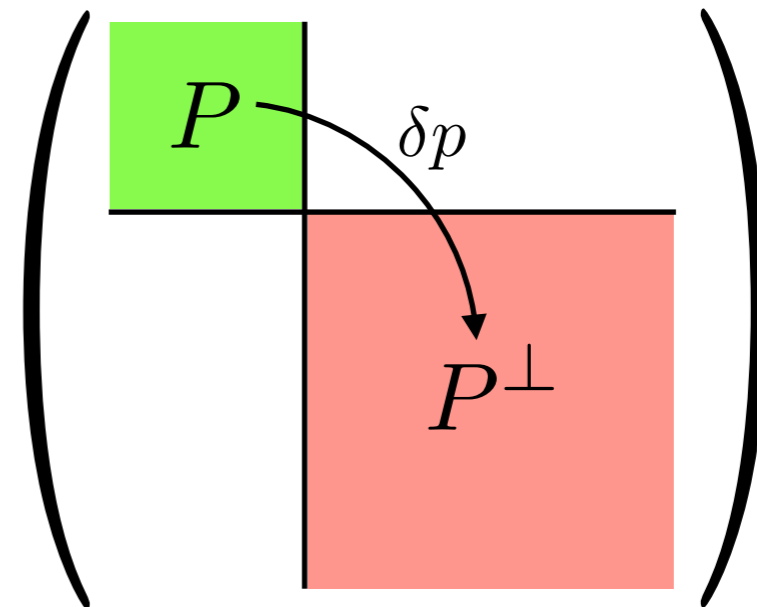
\nearrow jump ops

Generic RC Model



Error States

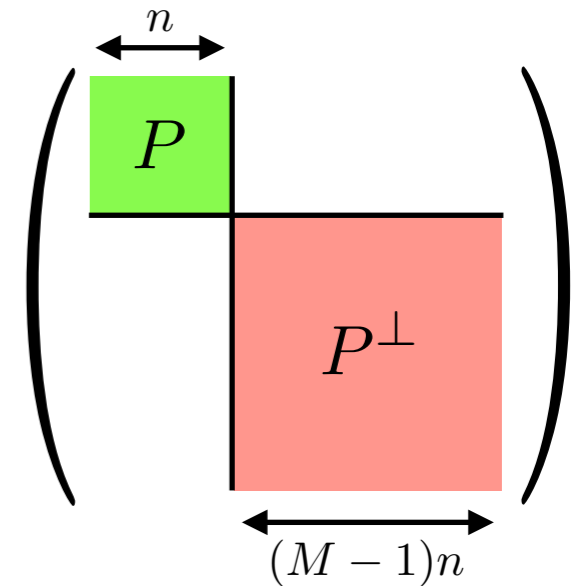
- projectors P P^\perp
- probability of error δp
- after small time δt



$$\begin{aligned} \delta p &= \text{tr}[P^\perp (\rho + \delta t \dot{\rho} + \frac{1}{2} \delta t^2 \ddot{\rho})] + \mathcal{O}(\delta t^3) \\ &= \text{tr}[P^\perp (\delta t \mathcal{L}\rho + \delta t^2 (H\rho H + \frac{1}{2} i[\mathcal{L}\rho, H] + \frac{1}{2} \frac{d}{dt} \mathcal{L}\rho))] \end{aligned}$$

Entropy Production

- pre-error $\rho = \rho P \quad S = S_0$
- post-error $S = S_0 + \log(M - 1)$
- entropy production



$$\begin{aligned} \delta\Sigma &= [-\delta p \log \delta p - (1 - \delta p) \log(1 - \delta p) + \\ &\quad \delta p(S_0 + \log(M - 1)) + (1 - \delta p)S_0] - [S_0] \\ &= \delta p(1 + \log(M - 1) - \log(\delta p)) + \mathcal{O}(\delta p^2) \\ &\geq \delta p \log M \end{aligned}$$

Quantum Zeno RC

- weak Lindbladian $\gamma_i \lesssim \mathcal{O}(\delta t)$
- error probability

$$\begin{aligned} \delta p &= \delta t^2 \operatorname{tr}[\rho(H P^\perp H + \sum_i (\gamma_i / \delta t) L_i^\dagger P^\perp L_i)] + \mathcal{O}(\delta t^3) \\ &= \delta t^2 \operatorname{tr}[\rho_{00}(U_{01} U_{01}^\dagger + \sum_i (\gamma_i / \delta t) L_{i,01}^\dagger L_{i,01})] \\ &= \delta t^2 \langle U' U'^\dagger \rangle \end{aligned}$$

effective thermal
potential

- QZE $\delta p \sim \mathcal{O}(\delta t^2)$

Quantum Zeno RC

- entropy production rate

$$\dot{\Sigma} \geq \delta t \langle U' U'^{\dagger} \rangle m \log g$$

- effective potential couples to each thermal dof nm

- coupling uncorrelated $\langle U' U'^{\dagger} \rangle \approx nmT^2$

effective
temperature



$$\dot{\Sigma} \geq \delta t nm^2 T^2 \log g$$

Quantum Zeno RC

- Zeno and computation rates (M-L)

$$\begin{aligned} \nu_Z &= 1/\delta t & \nu_C &\leq E_C/\pi n \\ &\leq E_Z/\pi & &= nm\varepsilon/\pi n \end{aligned}$$

- entropy generation rate // computation rate

$$\dot{\Sigma} \geq \frac{\nu_C^2}{\nu_Z} \underbrace{\frac{n\pi^2 T^2}{\varepsilon^2}}_{\eta} \log g \quad (\Delta\Sigma = \dot{\Sigma}/\nu_C)$$

- adiabatic bound!

$$\Delta\Sigma = \nu_C(\eta/\nu_Z) \propto \nu_C \quad \nu_C \leq \sqrt{\frac{1}{\pi\eta} \dot{\Sigma} E}$$

Strong Lindbladian RC

- strong Lindbladian

$$\gamma_i \sim \mathcal{O}(1)$$

- lose QZE

$$\delta p = \delta t \operatorname{tr}[\rho \sum_i \gamma_i L_i^\dagger P^\perp L_i]$$

- retain adiabaticity

$$\dot{\Sigma} \geq nm^2 T^2 \log g = \eta \nu_C^2$$

- BUT lose geometric scaling??

$$\nu_C \lesssim \sqrt{\dot{\Sigma}} \sim r$$

Strong Lindbladian RC

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- \Rightarrow subdivide

$$\nu_C \lesssim \sqrt{N \dot{\Sigma}} \sim r^{5/2}$$

Strong Lindbladian RC

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$$\nu_C \lesssim \sqrt{\dot{\Sigma}} \sim r$$

- \Rightarrow subdivide

$$\nu_C \lesssim \sqrt{N \dot{\Sigma}} \sim r^{5/2}$$

- (individually slow)

$$\nu_{C,1} \lesssim r^{-1/2}$$

physics of computing

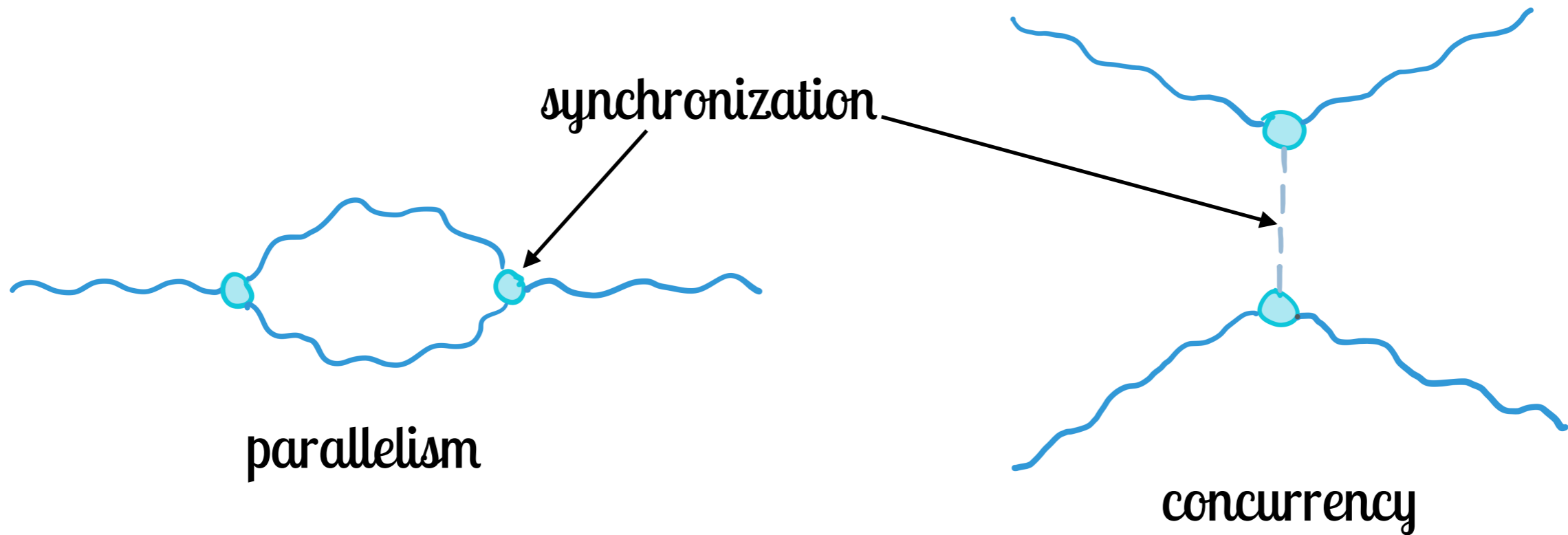
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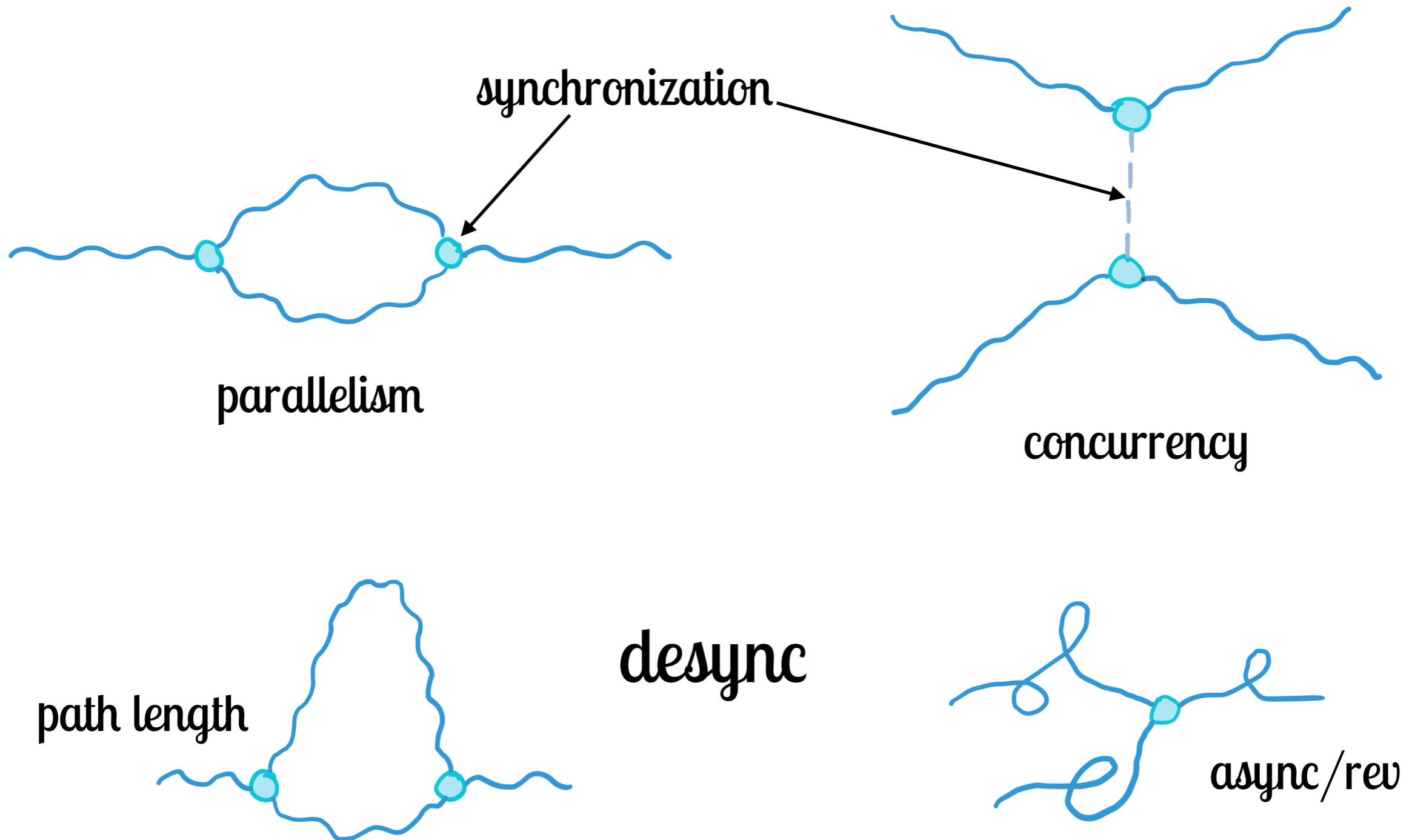
~~super-adiabatic RC?~~

parallelism & concurrency in RC

Synchronization



Synchronization



Async/Markovian RC

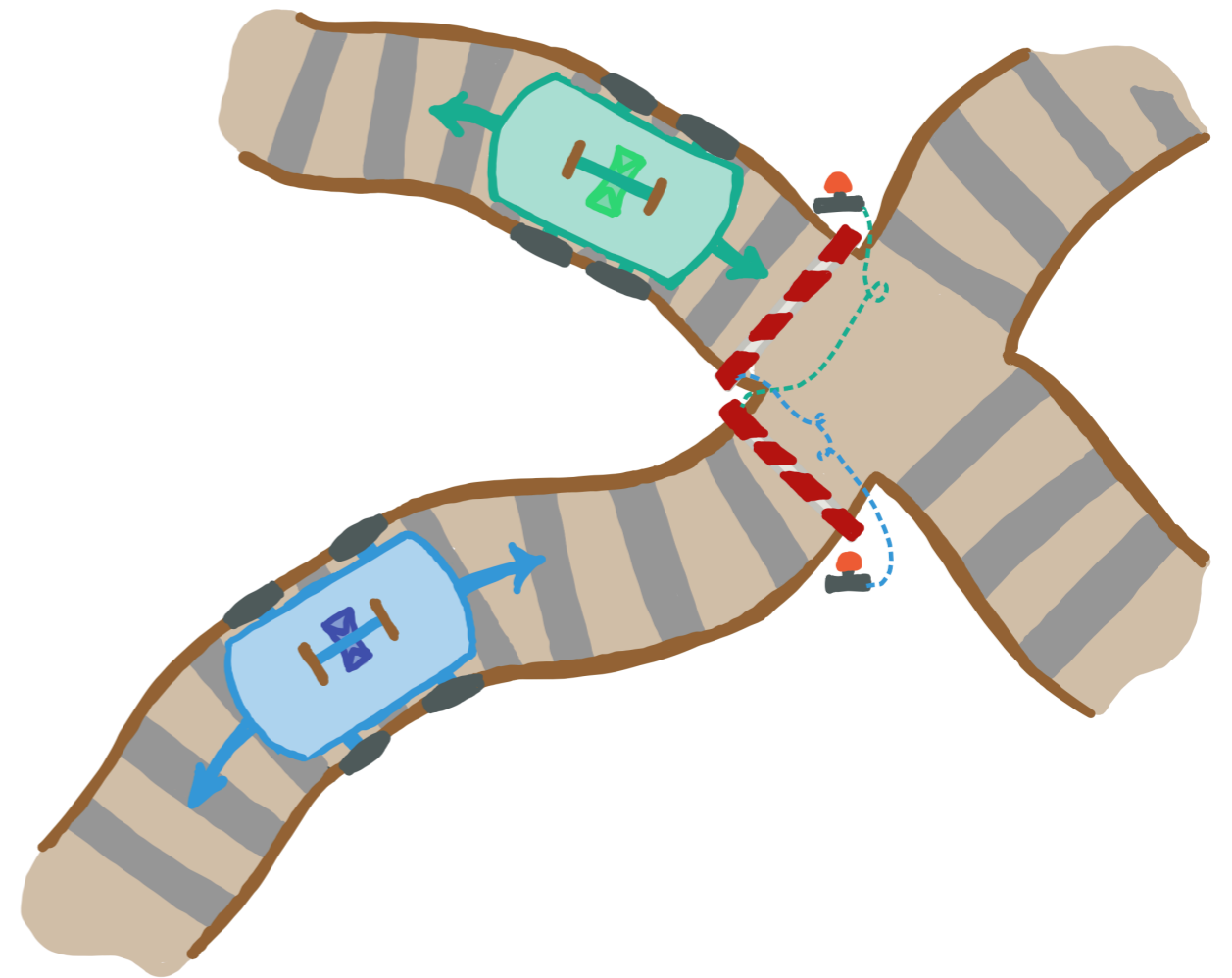
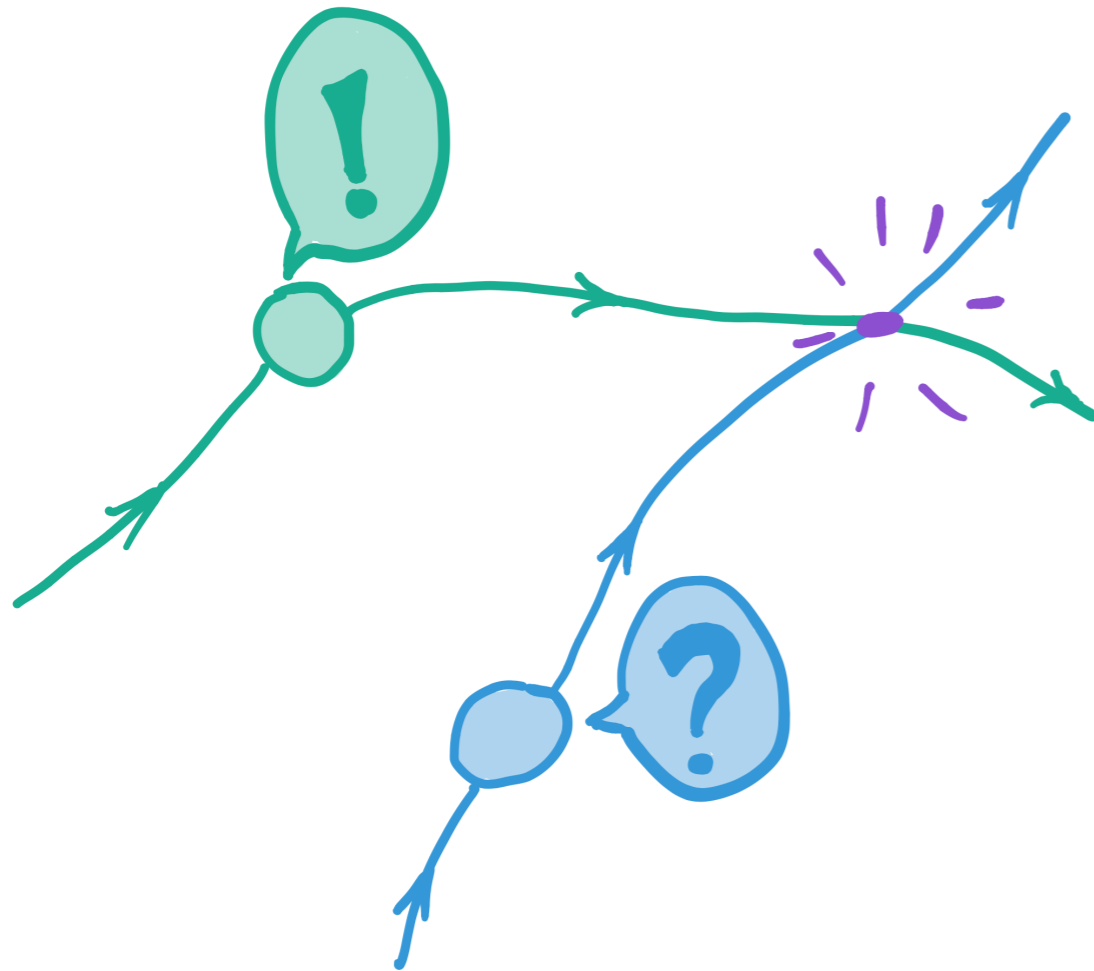


- realizes the adiabatic bound

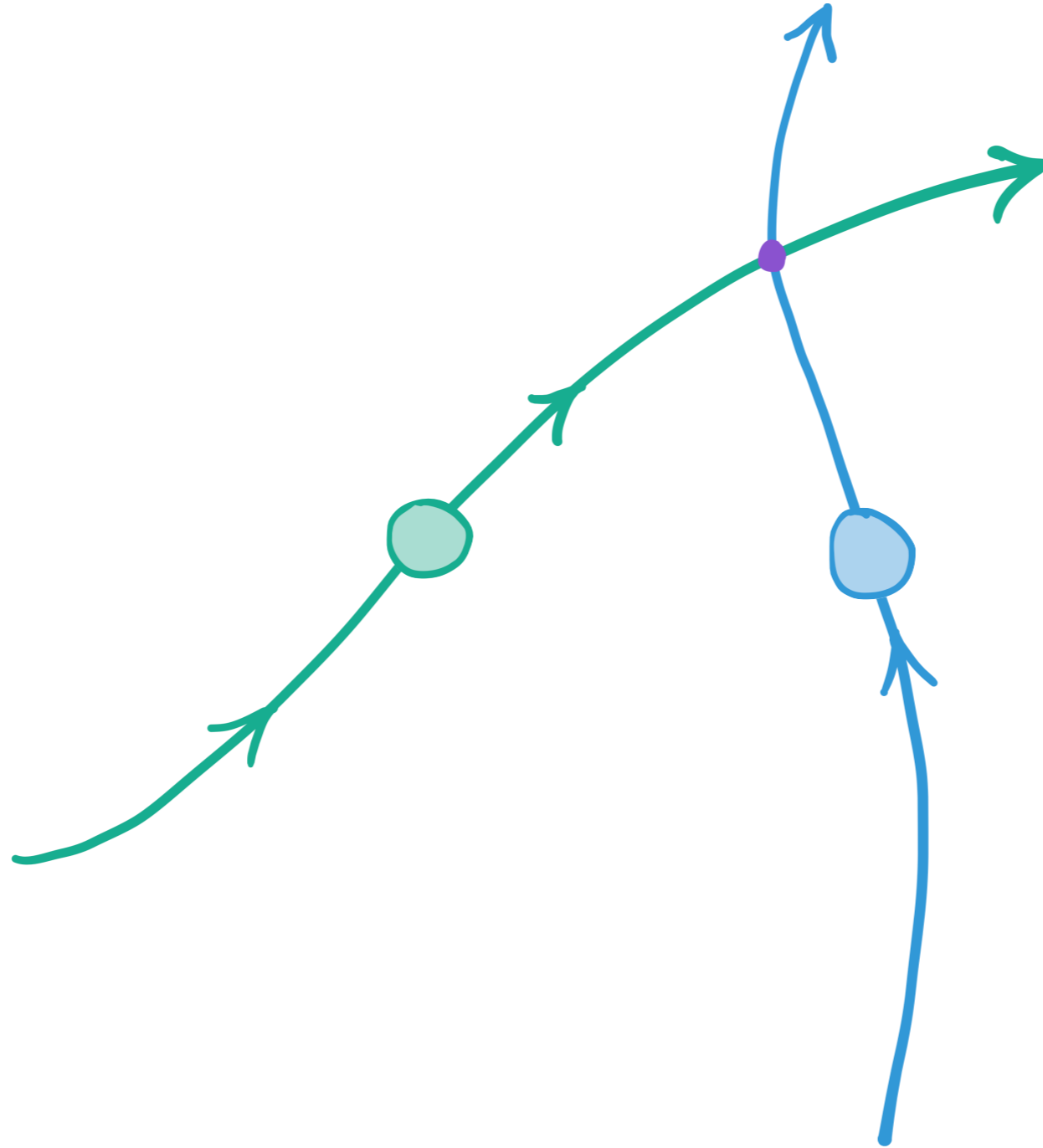
$$\nu_C \lesssim r^{5/2} \sim V^{5/6}$$

$$\nu_{C,1} \lesssim r^{-1/2} \sim V^{-1/6}$$

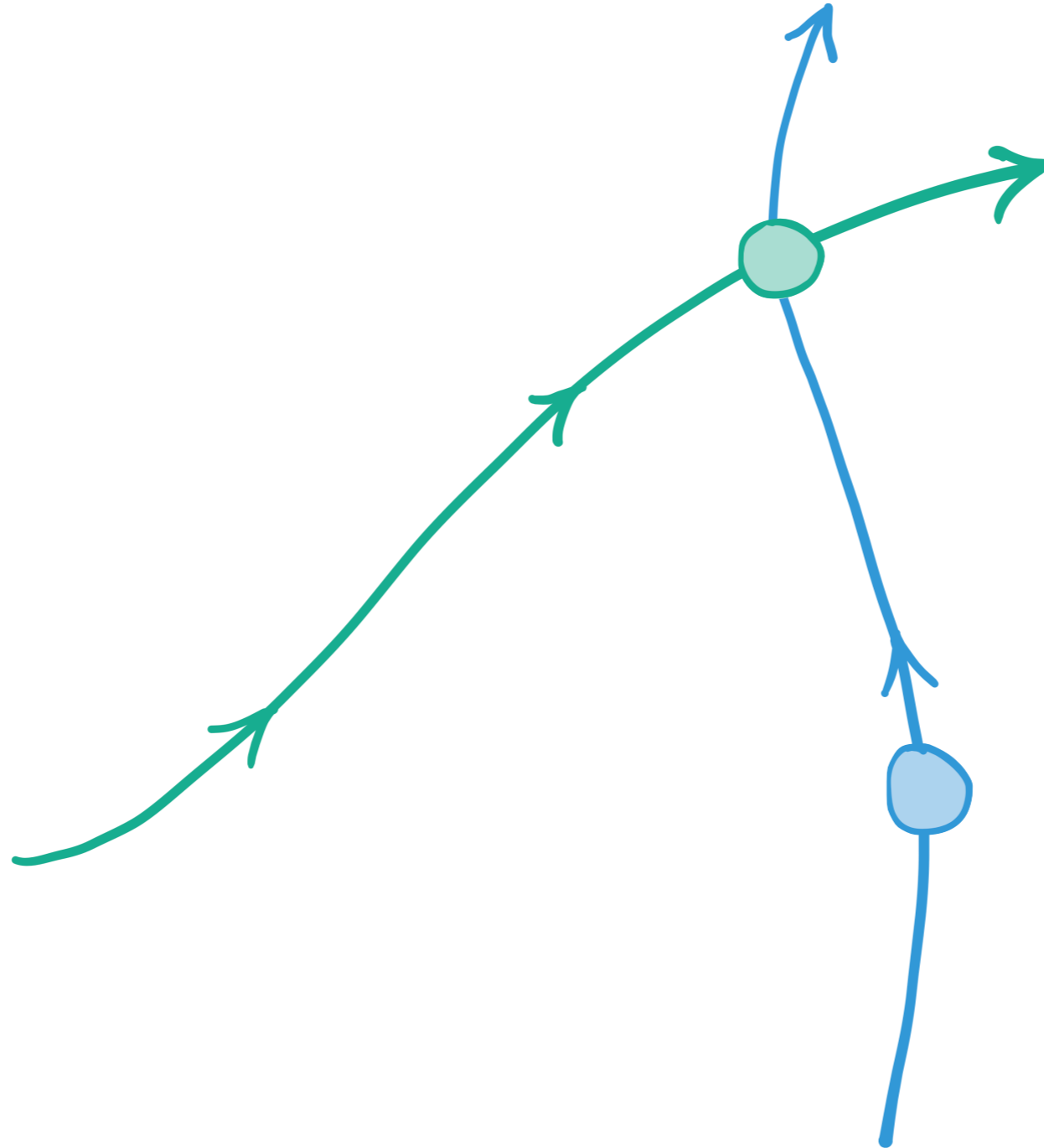
Asynchronous Sync



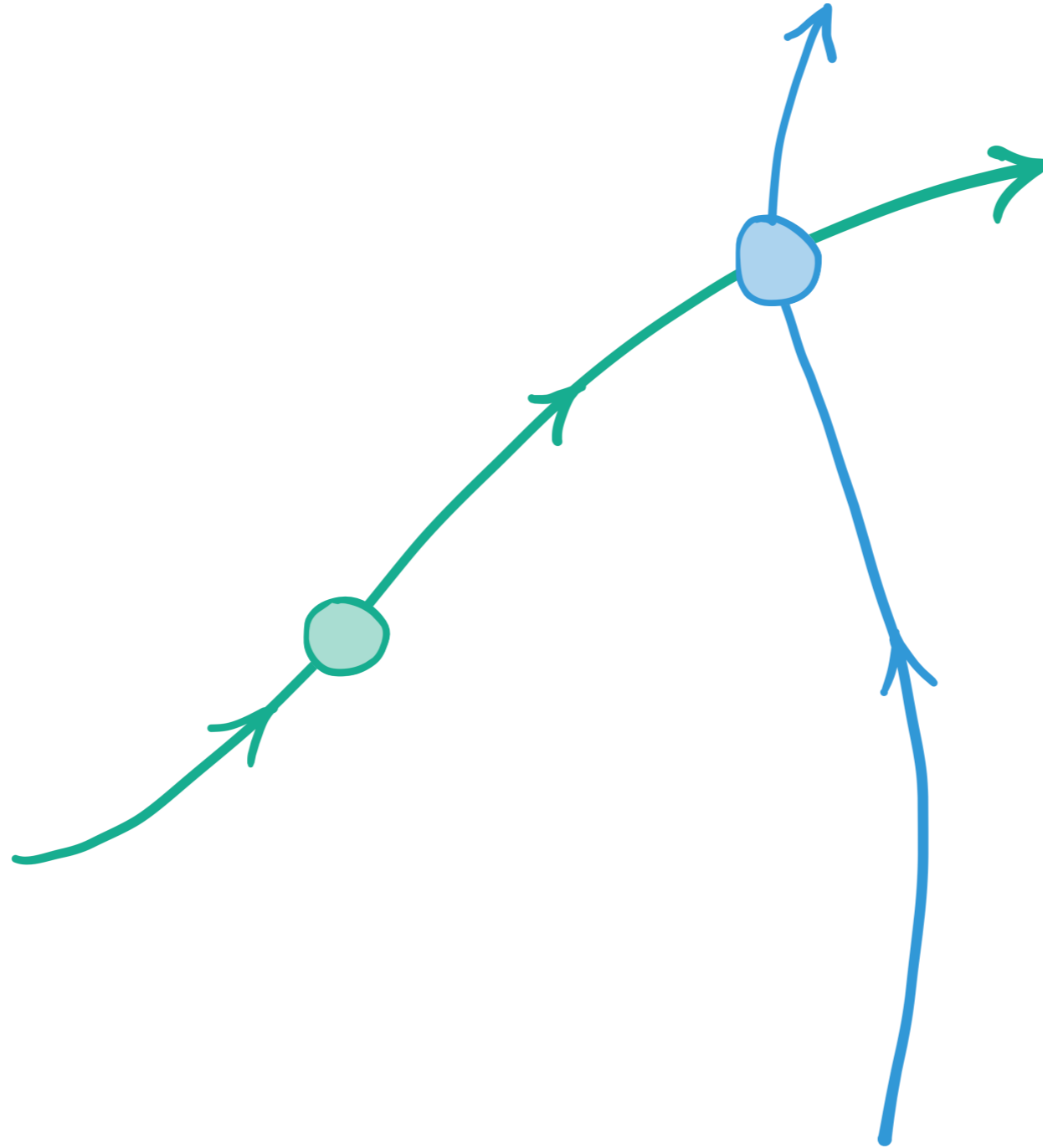
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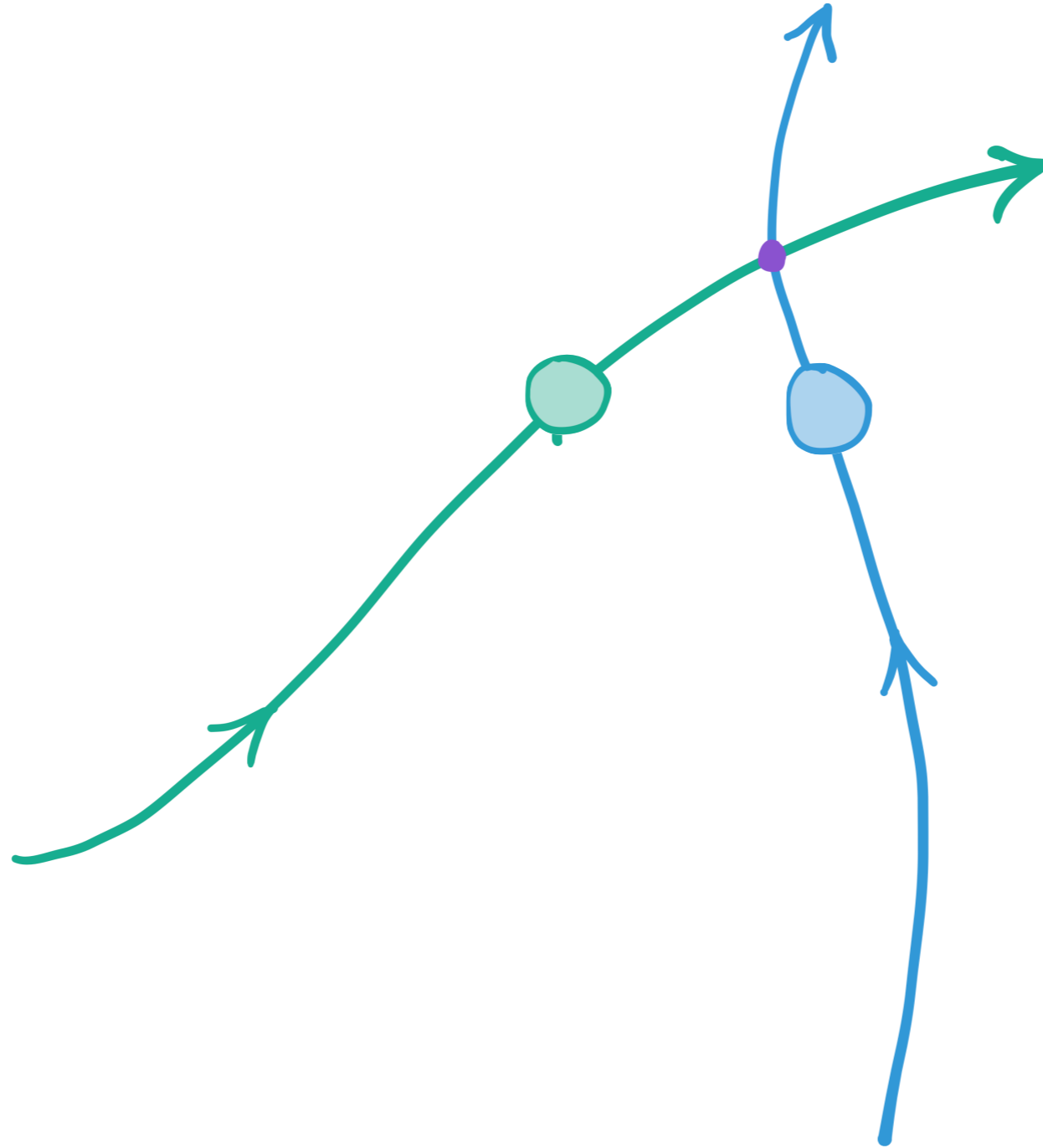
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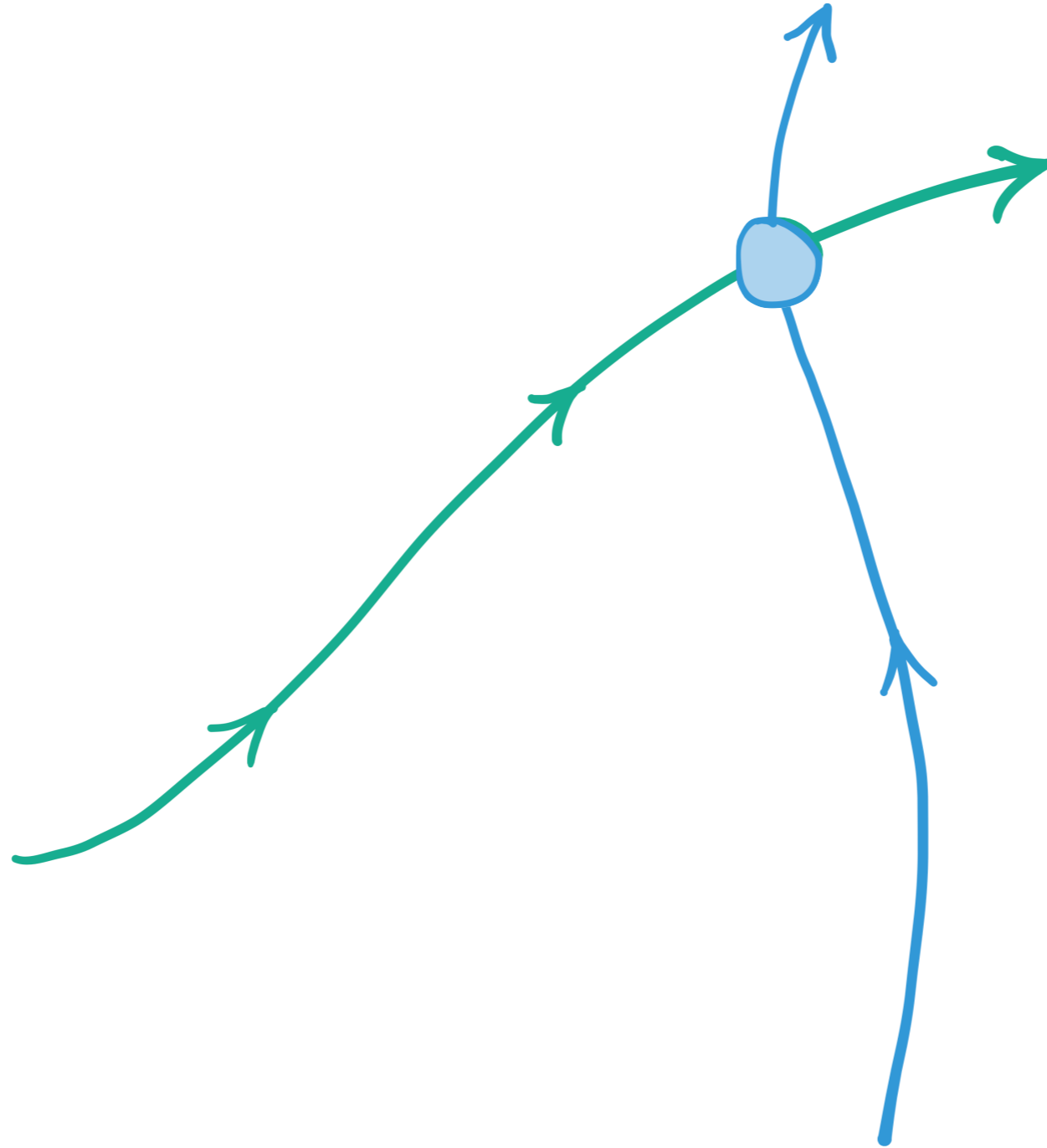
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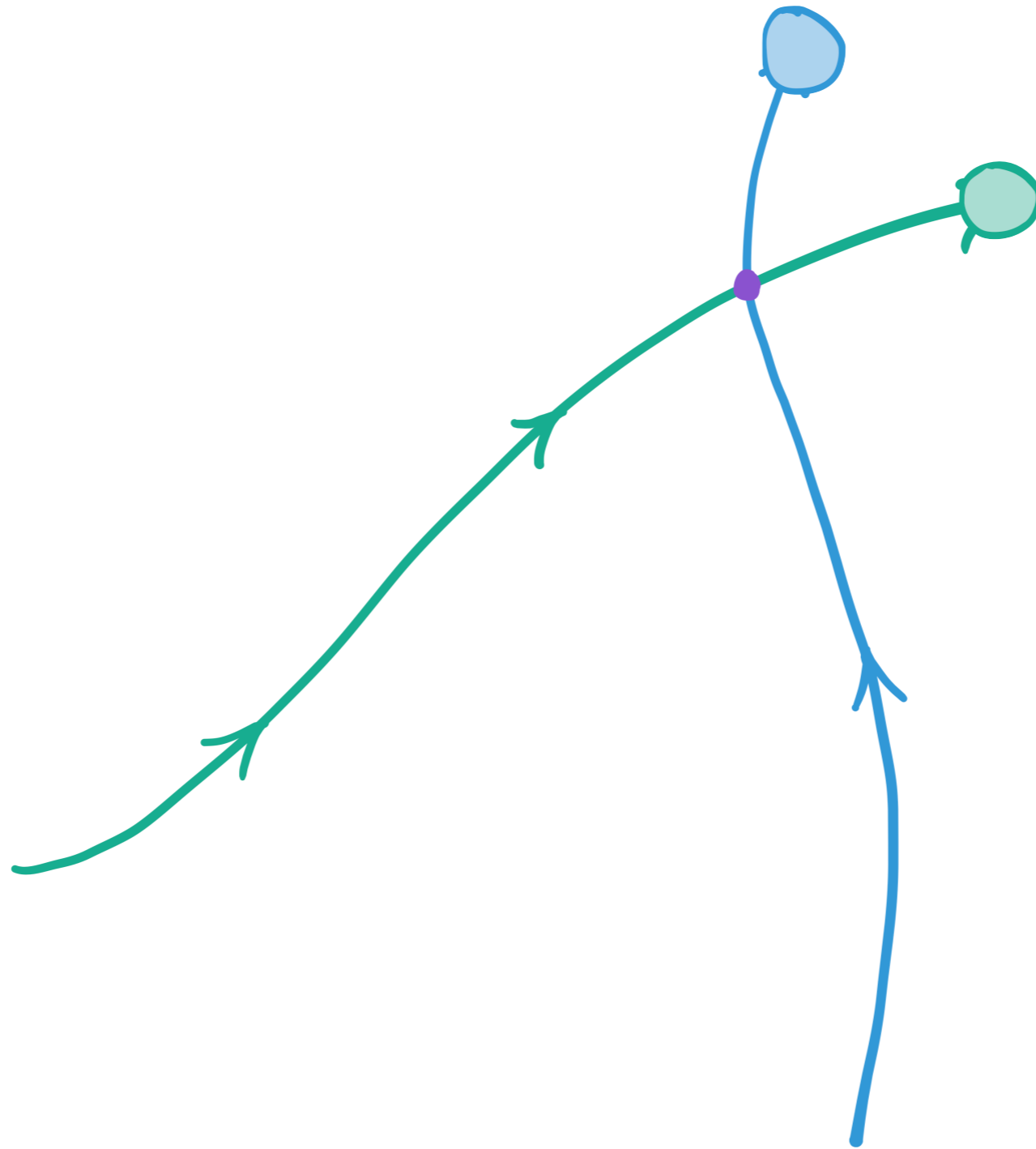
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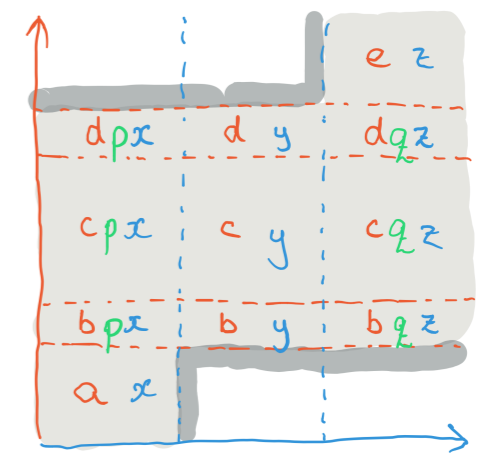
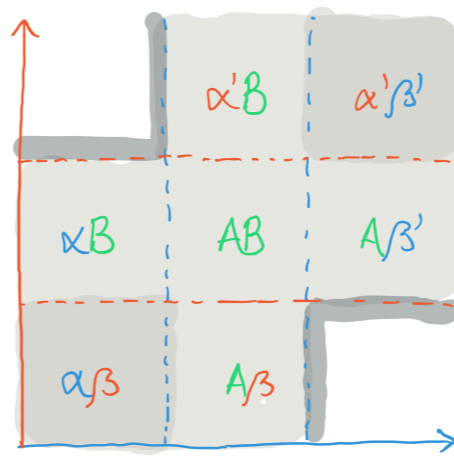
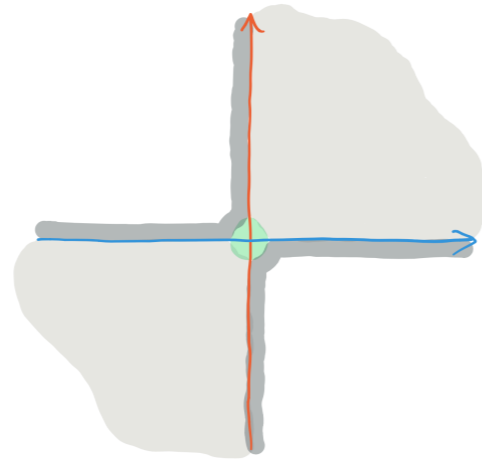
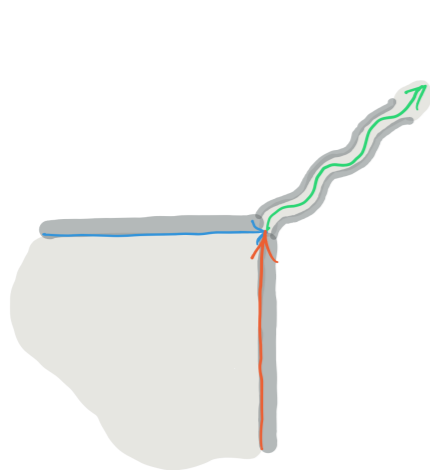
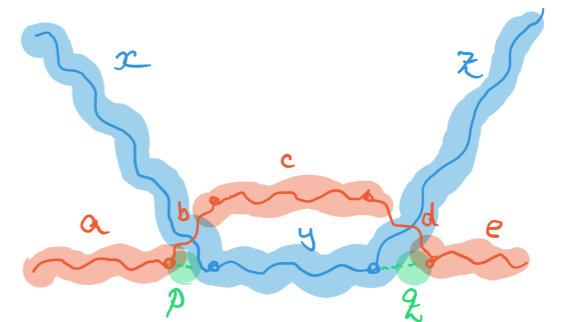
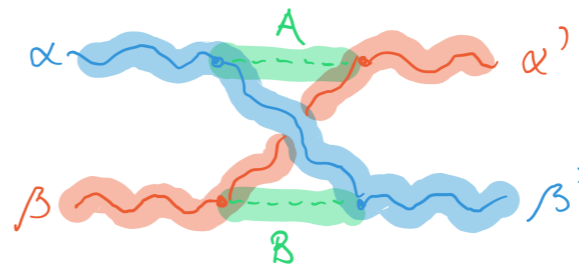
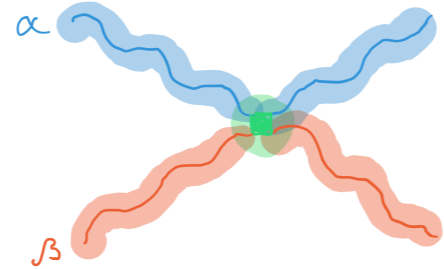
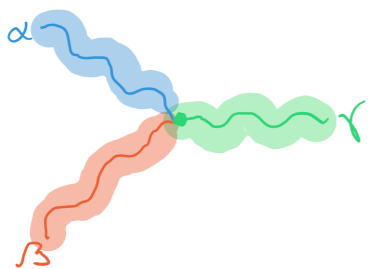
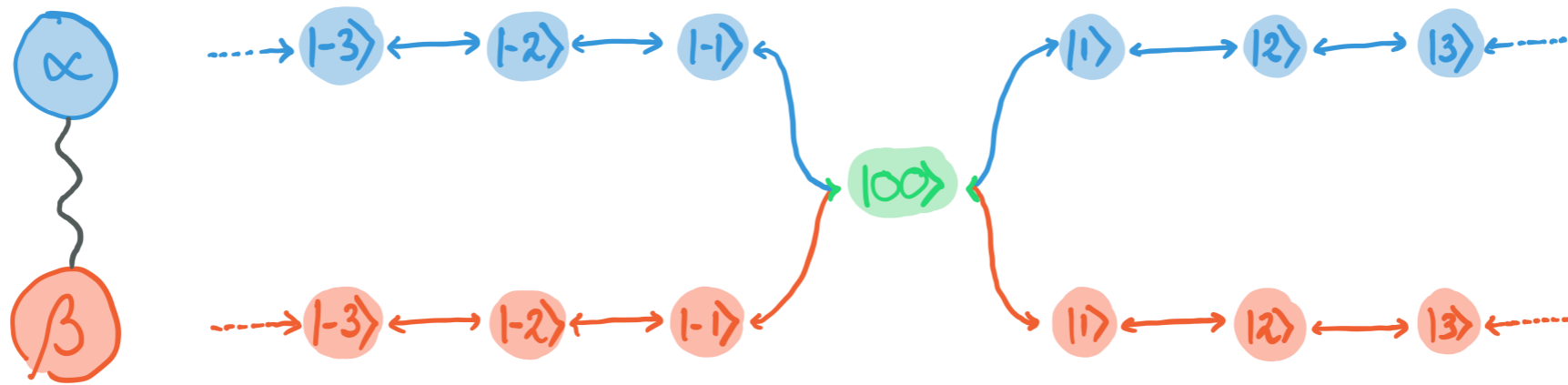
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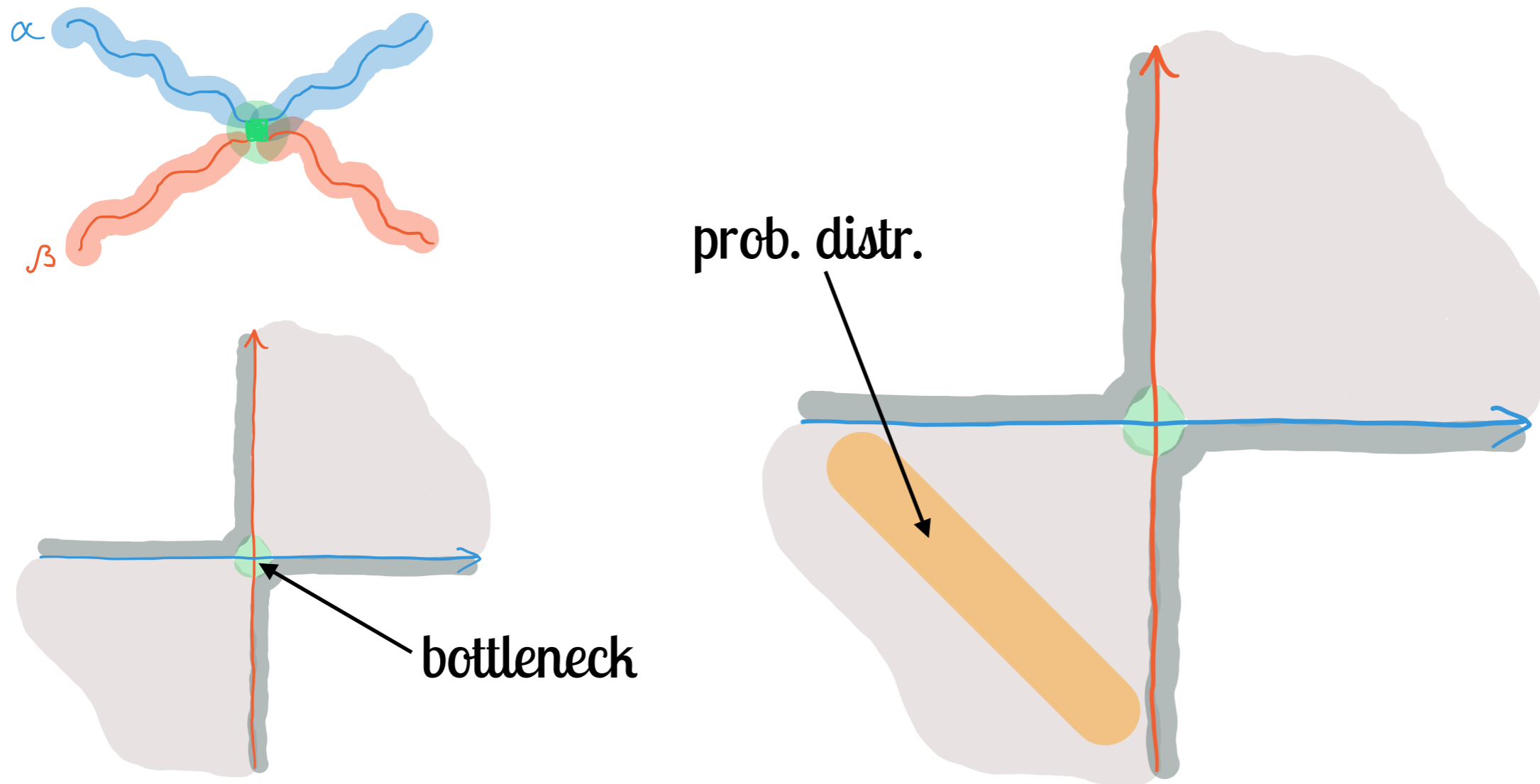
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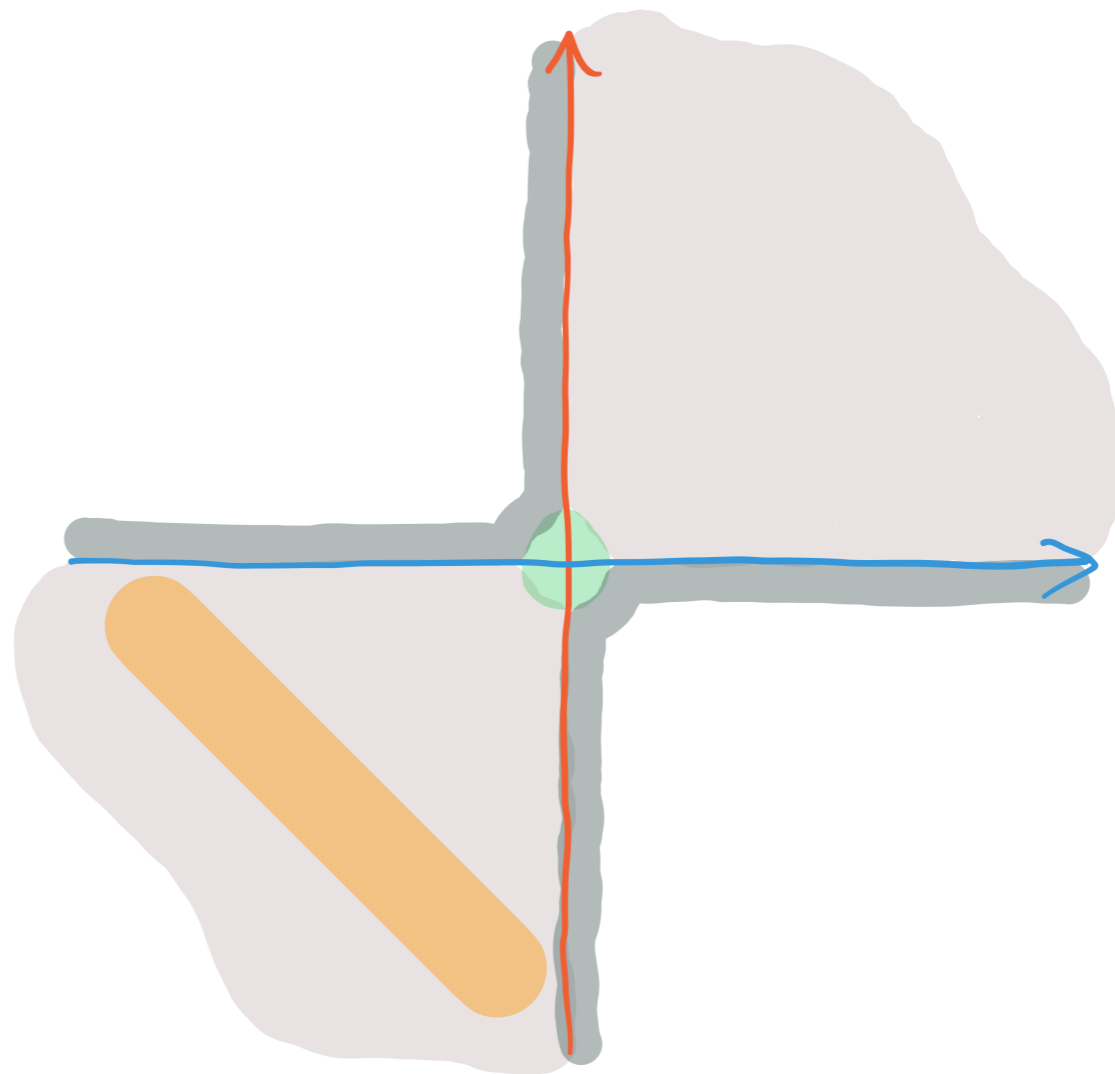
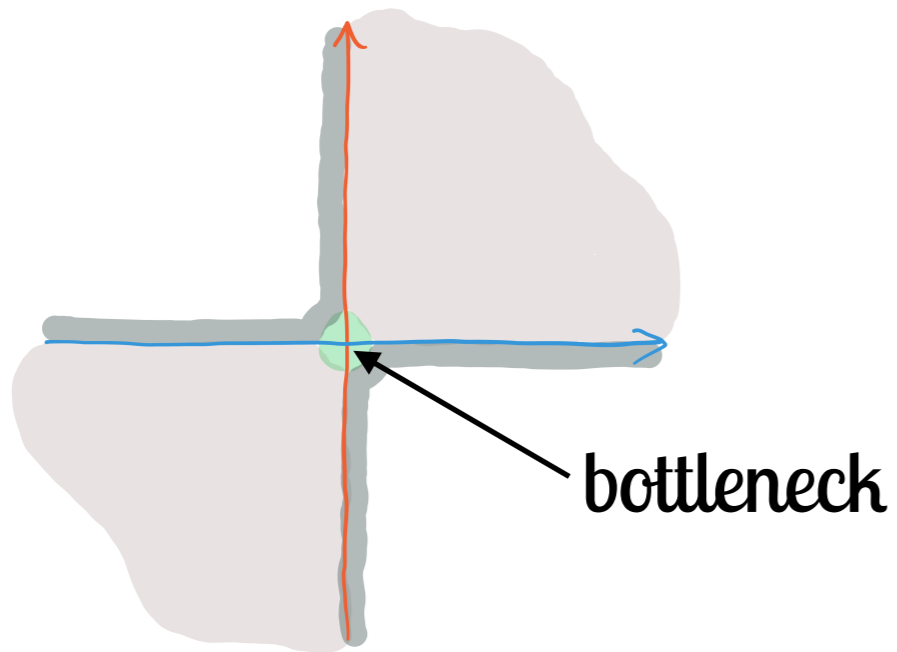
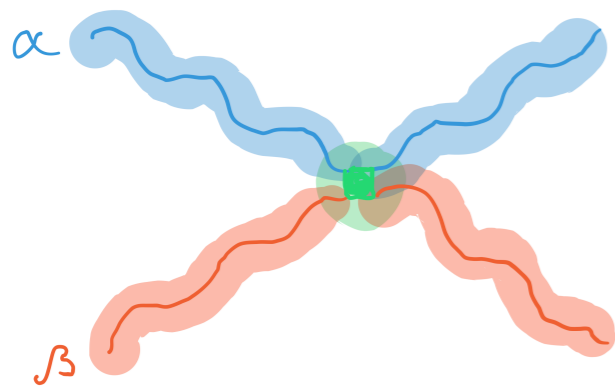
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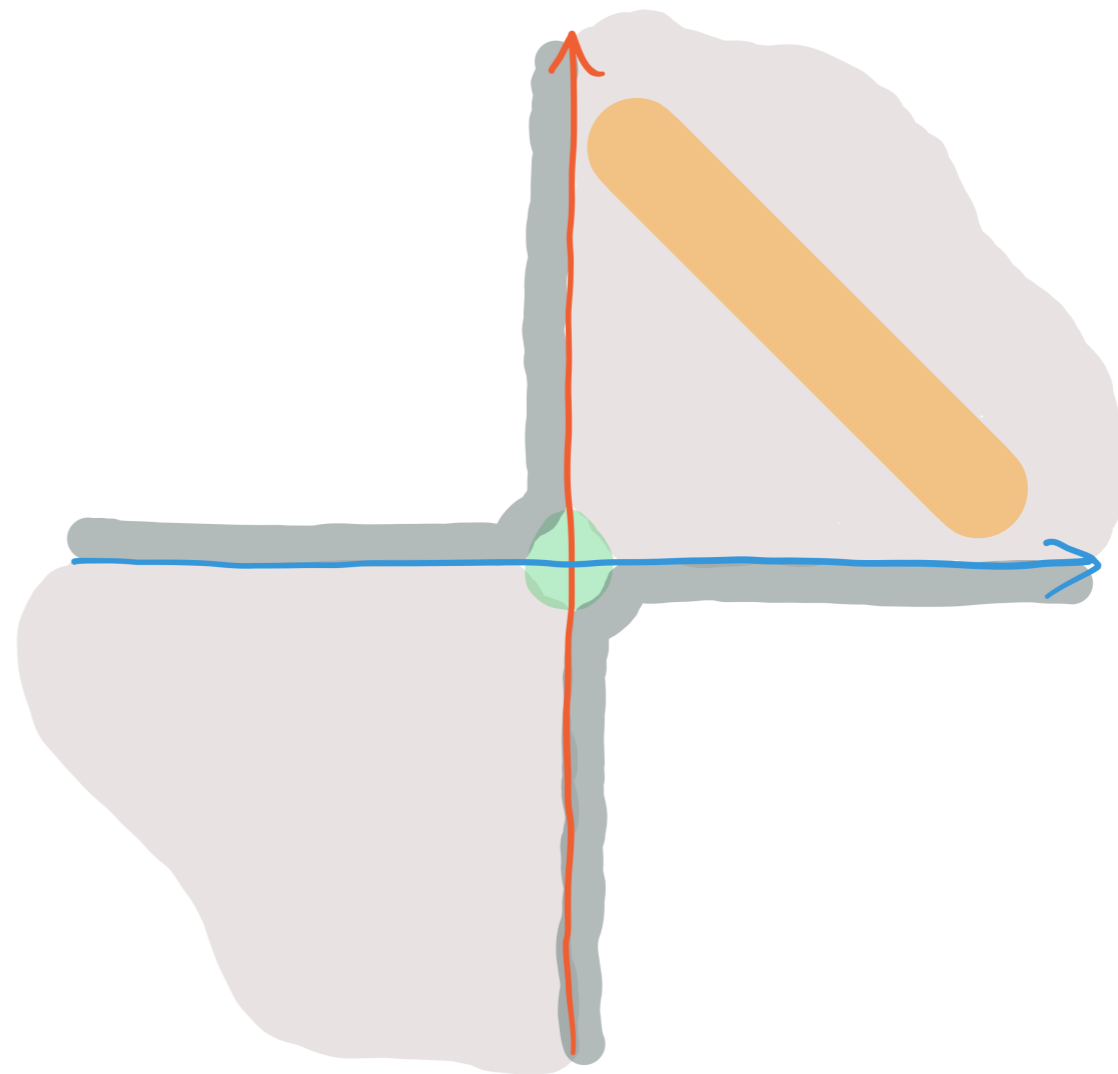
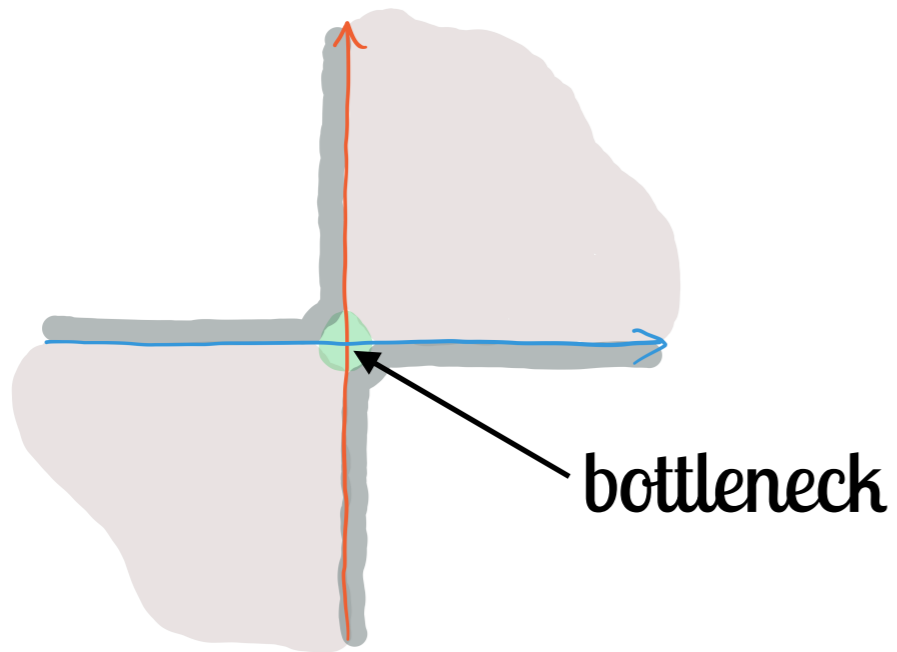
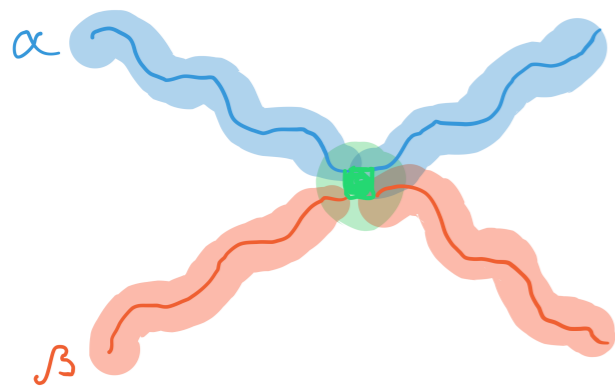
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Asynchronous Sync



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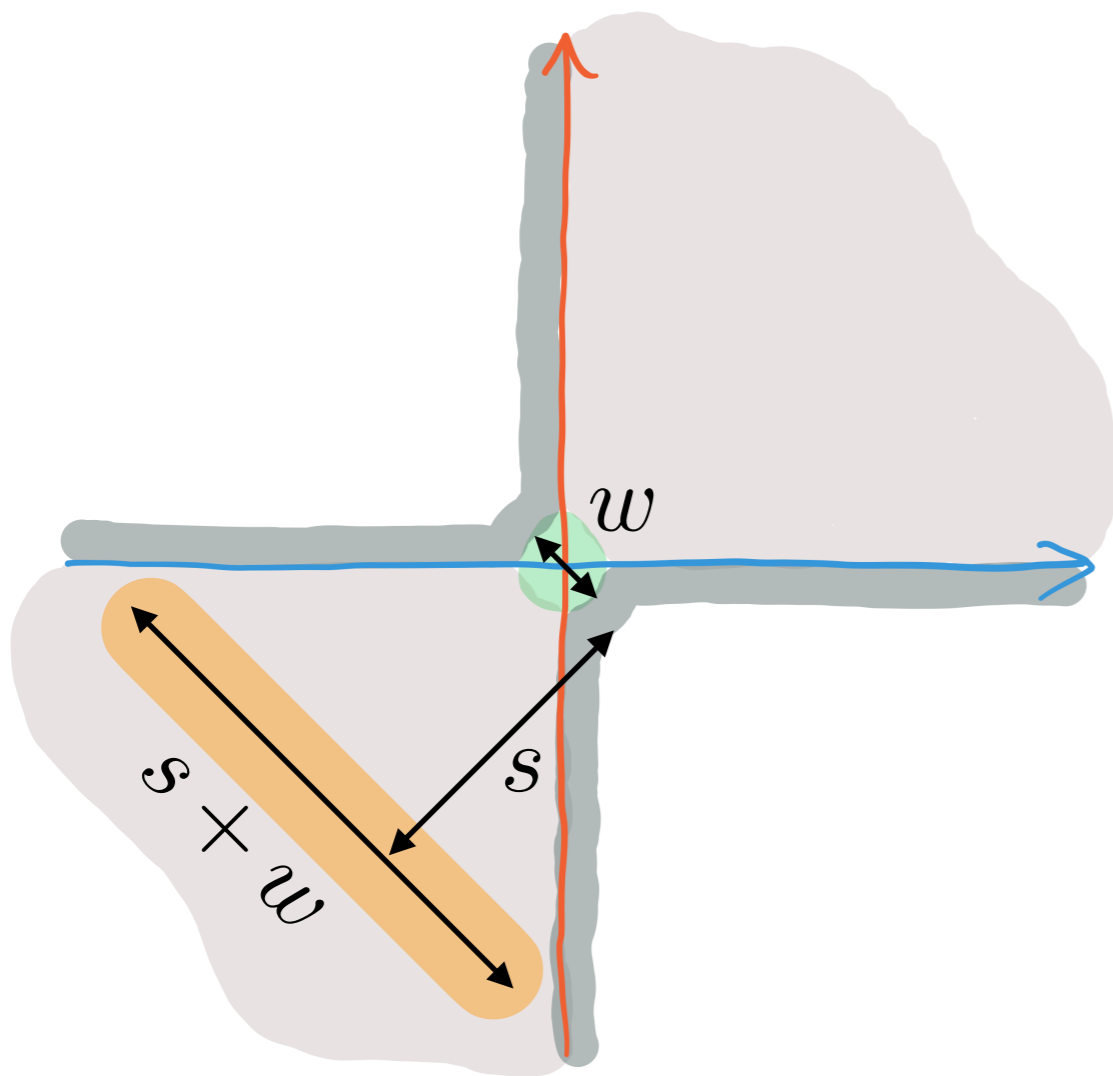
- info erasure model

$$\Delta I \approx \log \frac{s+w}{w}$$

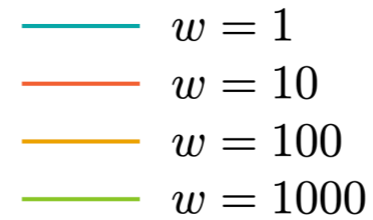
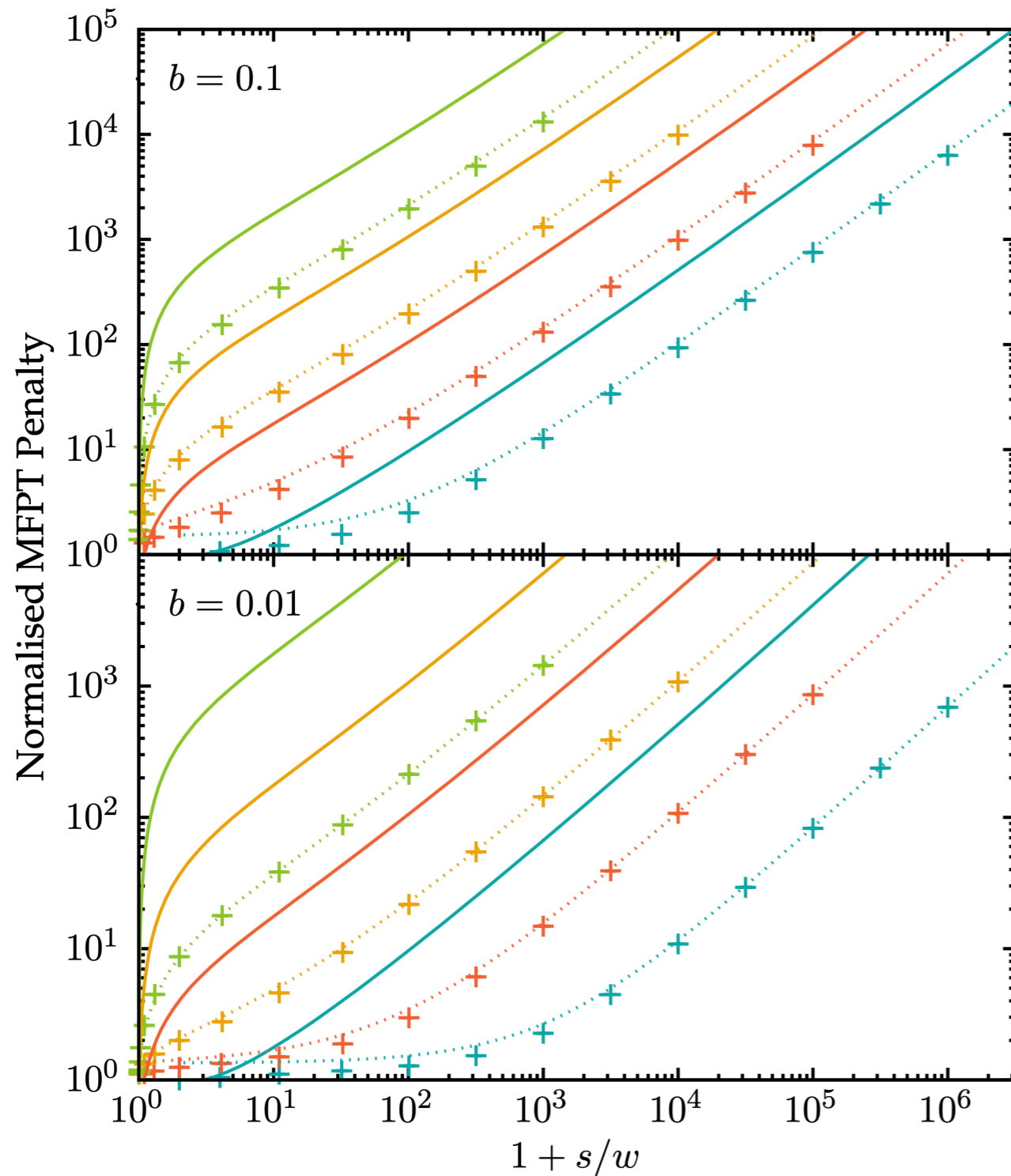
$$\Delta t \gtrsim \Delta I / \nu_{\text{irrev.},1}$$

$$\nu_{\text{irrev.},1} \sim r^{-1}$$

$$\nu_{\text{rev.},1} \sim r^{-1/2}$$



Asynchronous Sync



- numerical sim
- UB and LB from MFPT and discrete MC
- LB ~ info theory

Asynchronous Sync

- expensive, ~equivalent to entropy erasure
- significant parallelism/concurrency in async RC
⇒ comp rate falls to ~irrev
- constrains algorithms to minimise sync or constrains net comp rate

Conclusions & Future Work

- RC > irrevC; adiabatic upper bound
 - for large & sustained comp
 - concrete const factors?
 - super-adiabatic special cases? (e.g. Pidaparthi+Lent)
- async sync expensive
 - sync sync performance?
 - preliminary results suggest cheap
 - algorithm design? mitigations for algorithmic desync?

Thank you!



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